

From previous lecture, we need to find the flows
 Find x that maximizes

$$L(x, \lambda) = f(x) - \sum \lambda_j g_j$$

$$= u(x_1) + u(x_2) + u(x_3) - \lambda_1 x_1 - \lambda_2 x_2 + \lambda_1 - \lambda_2 x_1 - \lambda_2 x_3 + \lambda_2$$

find flows

x_1
 x_2
 x_3 that maximize

$$u(x_1) - (\lambda_1 + \lambda_2) x_1$$

$$u(x_2) - \lambda_2 x_2$$

$$u(x_3) - \lambda_2 x_3$$

\Rightarrow x_1 adjusts flow independently of x_2 and x_3

Do this for a given value of λ obtaining

$$\underline{x}(\underline{\lambda})$$

then minimize with respect to λ i.e.

find shadow prices for links 1, 2

$$\lambda^* = \underset{\lambda}{\operatorname{argmin}} D(\lambda) = \underset{\lambda}{\operatorname{argmin}} \max_x L(x, \lambda)$$

This corresponds to finding the saddle point where the Lagrangian is minimum with respect to price and max with respect to flow.

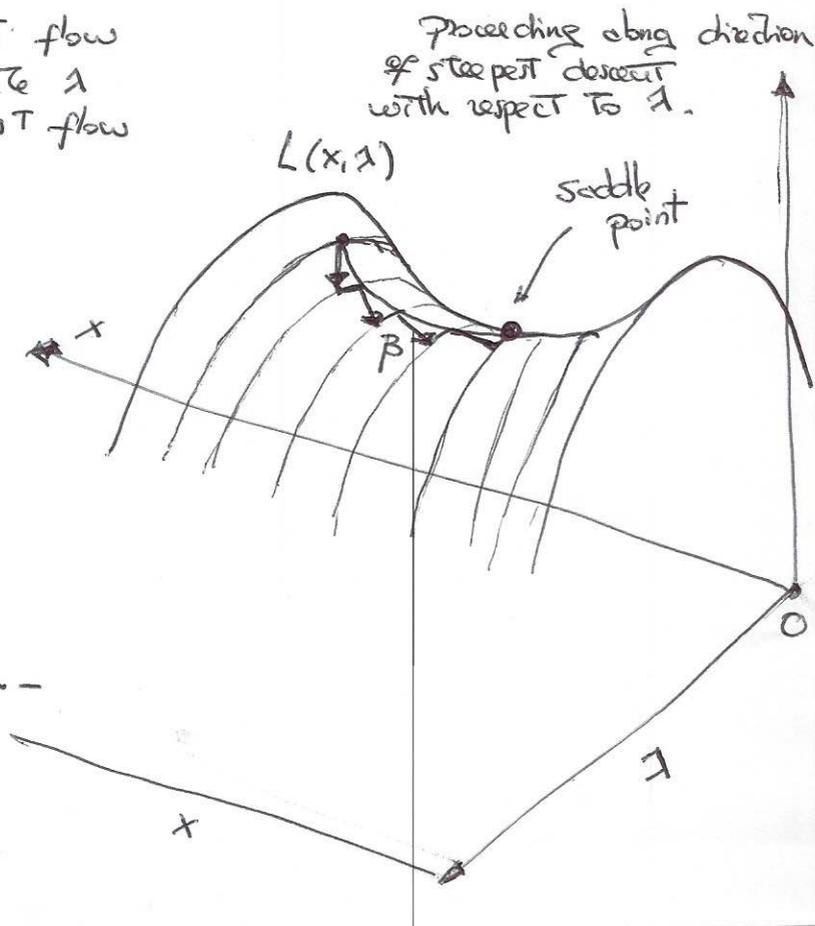
How can we do this distributively.

ITERATIVE ALGORITHM

adjust flow
 update λ
 adjust flow

$$\lambda_j(n+1) = \lambda_j(n) - \beta \frac{\partial L(x, \lambda)}{\partial \lambda_j}$$

This corresponds to taking a step of length β along the direction of steepest descent with respect to λ . The direction is tangential to the curve of steepest descent passing through the points where x is maximum.



At each step adjust the flows to be on the max point for the given $\lambda(n)$

$$\begin{cases} x_1(n) = \underset{x_1}{\operatorname{argmax}} u(x_1) - [\lambda_1(n) - \lambda_2(n)] x_1 \\ x_2(n) = \underset{x_2}{\operatorname{argmax}} u(x_2) - \lambda_2(n) x_2 \\ x_3(n) = \underset{x_3}{\operatorname{argmax}} u(x_3) - \lambda_2(n) x_3 \end{cases}$$

Then proceed along the line of steepest descent

$$\begin{cases} \lambda_1(n+1) = \lambda_1(n) - \beta \frac{\partial}{\partial \lambda_1} L(x, \lambda_1) \\ \lambda_2(n+1) = \lambda_2(n) - \beta \frac{\partial}{\partial \lambda_2} L(x, \lambda_2) \end{cases}$$

Repeat the steps

How are step 1 & step 2 achieved?

Look evolution of queue length for flow j over an interval of time τ

$$q_j[(n+1)\tau] = q_j[n\tau] + \tau g_j$$

where ~~g_j~~
 for example $g_1 = x_1 + x_2 - 1$
 = arrival rate - service rate -
 = excess rate = increment in queue length

This is remarkably similar to eq. for $\lambda(n)$ that we need.

So we set

$$\lambda_j' = \frac{\beta}{\tau} g_j$$

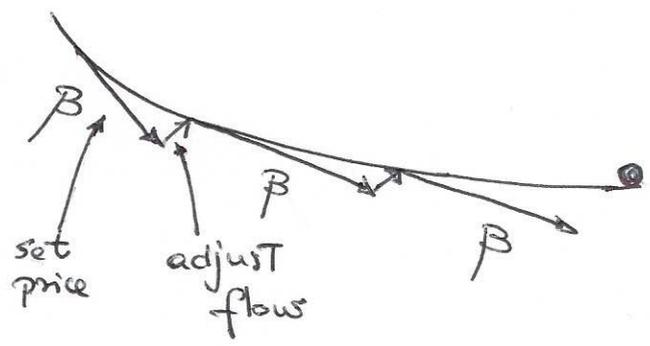
To have 1 step of gradient descent algorithm be executed every τ time units of queue evolution -

this either increases the price if $g_j > 0$ which forces the flow to decrease (via AIMD)

or decreases the price if $g_j < 0$ forcing the flow to increase -

We only need knowledge of queue length -
 With it we can set prices and then let flow be readjusted
 via AIMD -

Setting prices corresponds to taking a step towards saddle point -
 Adjusting flow corresponds to climbing up to maximum with
 respect to x for the given value of λ



Queue length is estimated via RED
 measuring packet losses at host that are proportional to
 queue length because router drops packets with probability
 proportional to queue length.