

Point-to-point codes for interference channels: A journey toward high performance at low complexity

Young-Han Kim

University of California, San Diego

Communication Theory Workshop
Dana Point, California
Tuesday, May 12, 2015

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SPECIAL
Report

WORDS BENNETT RING ILLUSTRATION NIGEL BUCHANAN

CONNECTIVITY

-IS KING-



From smart watches that synchronise with smartphones, to portable high-definition cameras that can be remotely monitored from anywhere on Earth, 2014 has been a year in which to be linked in is everything.

Where is wireless going?

Exabytes per Month

61% CAGR 2013-2018

18

9

0

2013

2014

2015

2016

2017

2018

- Mobile File Sharing (2.9%)
- Mobile M2M (5.7%)
- Mobile Audio (10.6%)
- Mobile Web/Data (11.7%)
- Mobile Video (69.1%)

Figures in parentheses refer to traffic share in 2018.

Source: Cisco VNI Mobile, 2014

Where is wireless going?

Billions of Devices

8% CAGR 2013–2018

12

6

0

2013

2014

2015

2016

2017

2018

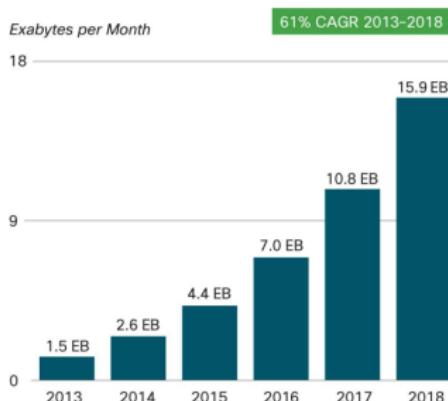
- Other Portable Devices (0.3%, 0.3%)
- Tablets (1.3%, 5.0%)
- Laptops (2.1%, 2.6%)
- M2M (4.9%, 19.7%)
- Smartphones (24.9%, 38.5%)
- Non-Smartphones (66.4%, 33.9%)

Figures in parentheses refer to device or connections share in 2013, 2018.

Source: Cisco VNI Mobile, 2014

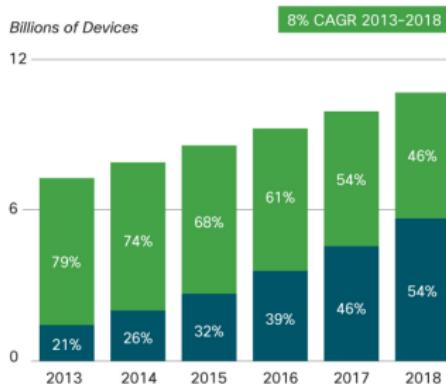
Where is wireless going?

Mobile data per month



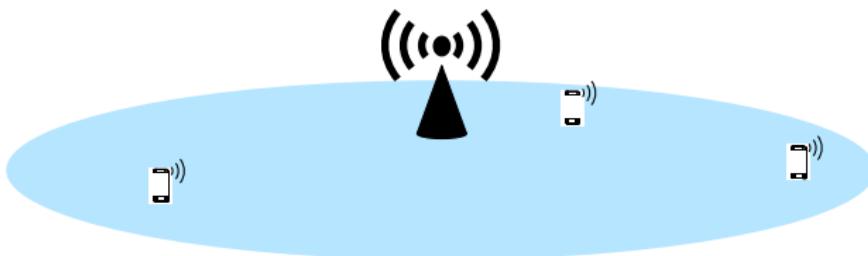
Source: Cisco VNI Mobile, 2014

Number of devices



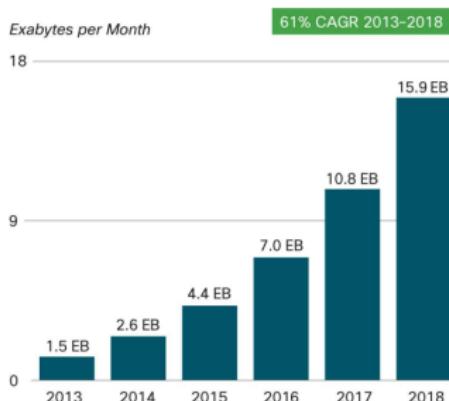
Percentages refer to device or connections share.

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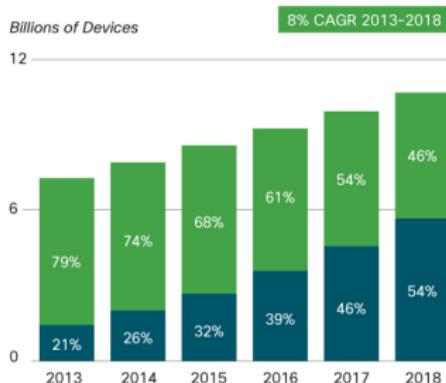
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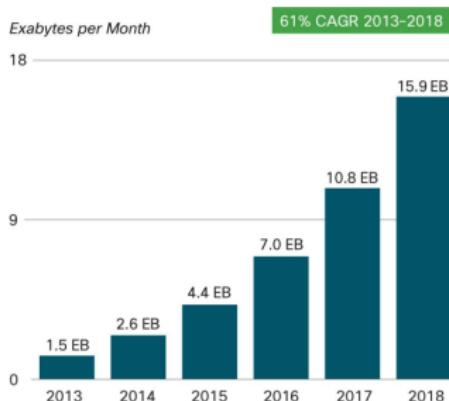
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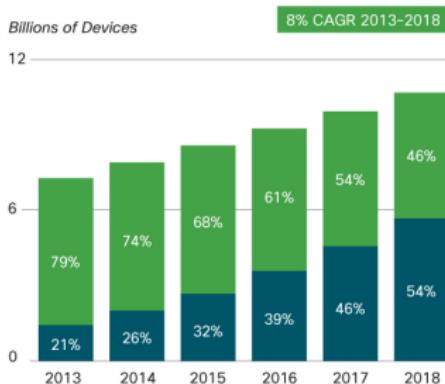
Where is wireless going?

Mobile data per month



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Number of devices



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The Internet of THINGS

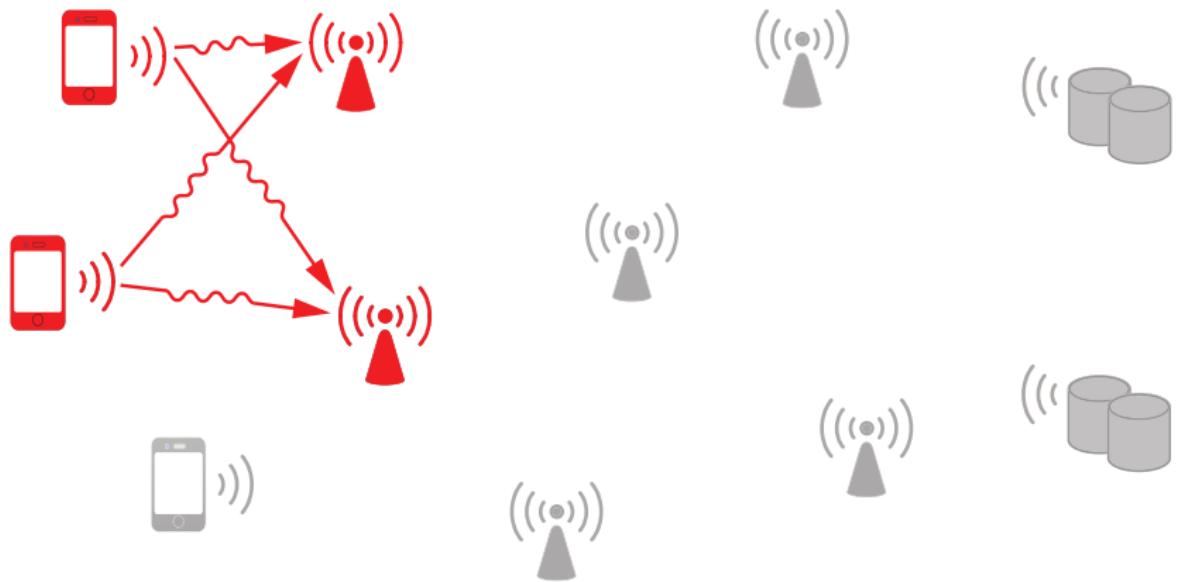


CONNECT
THE WORLD

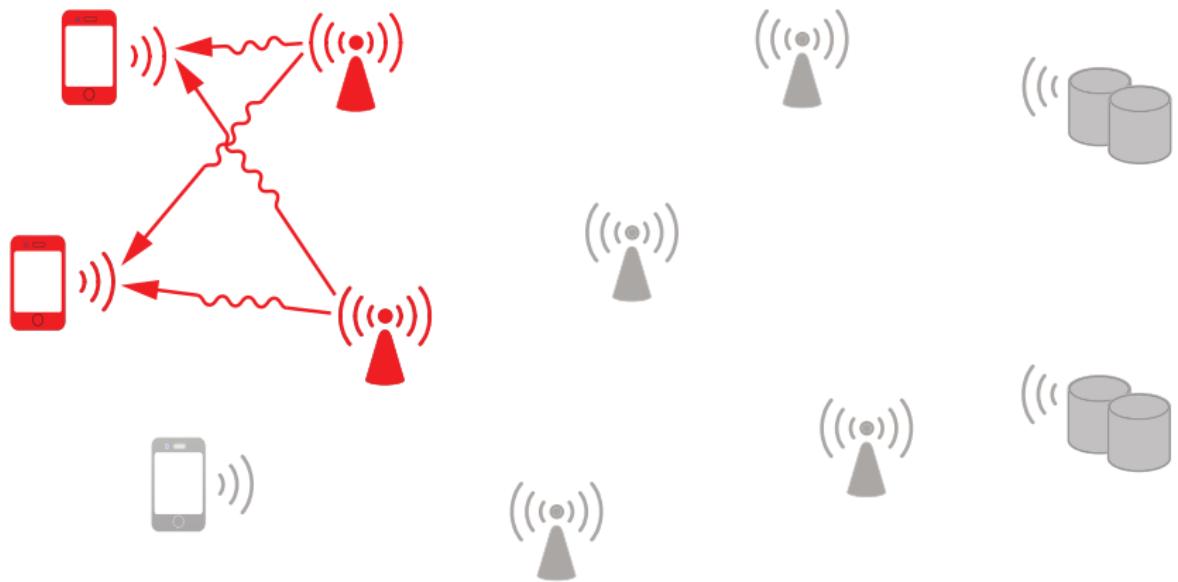
Interference, interference, interference



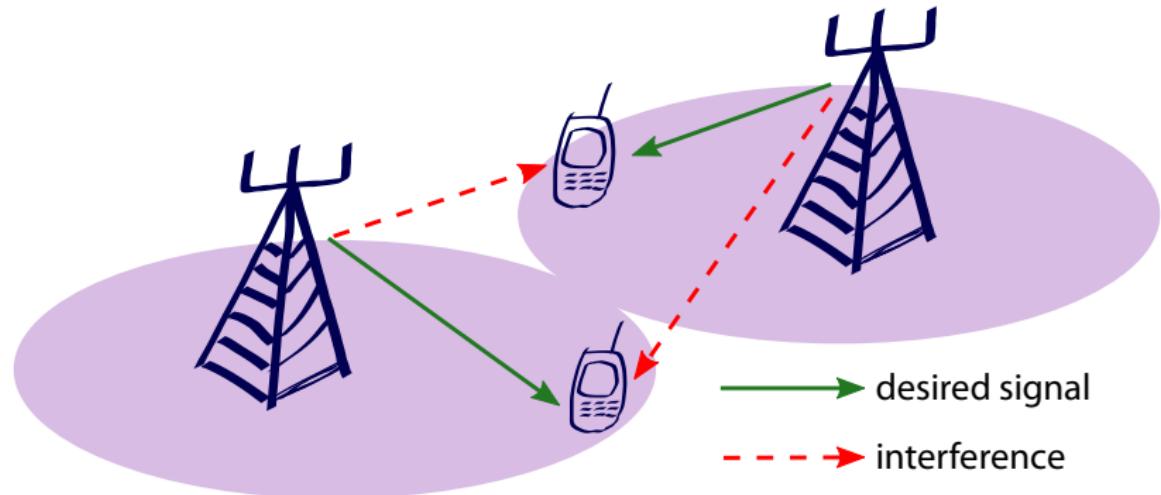
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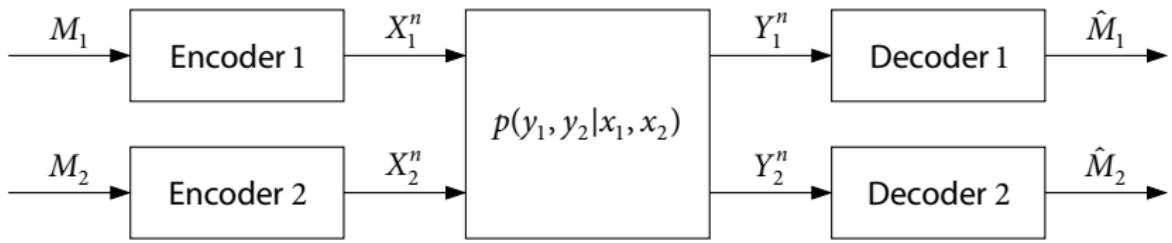
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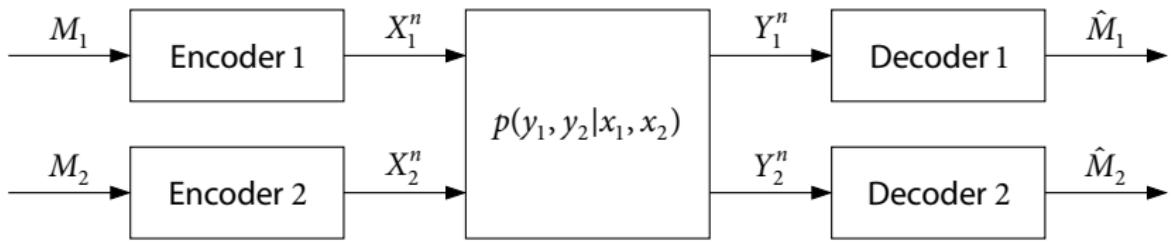
Interference in cellular networks



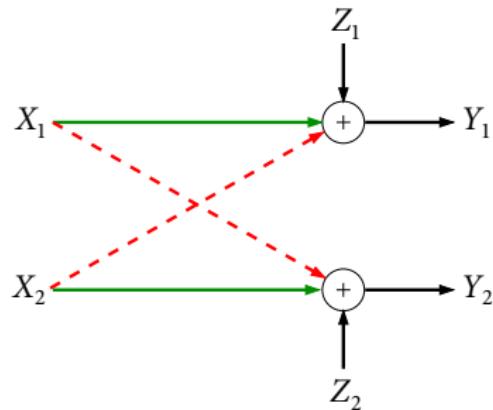
Interference channel



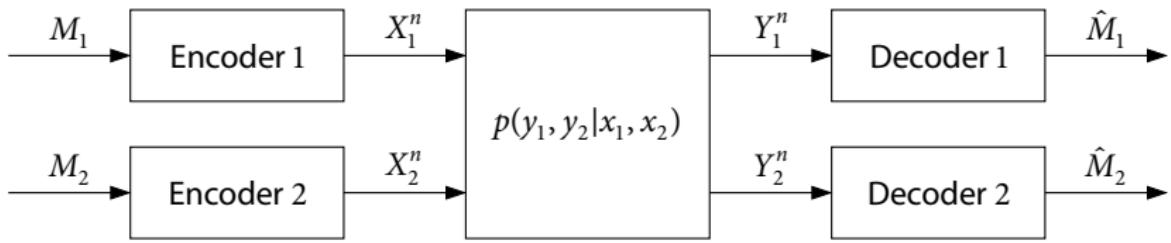
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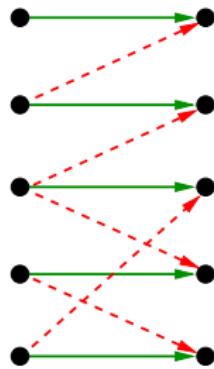
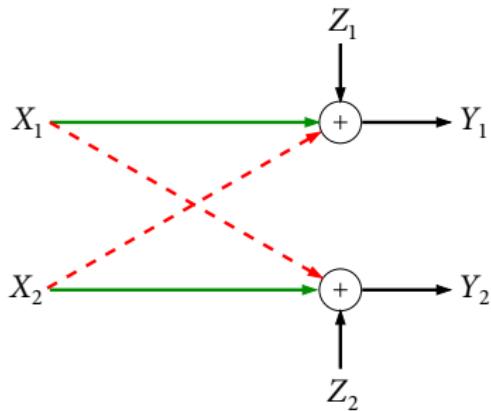
- Gaussian interference channel



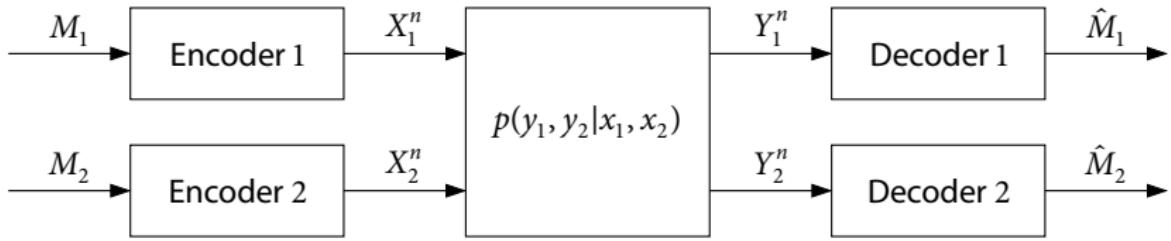
Interference channel



- Gaussian interference channel
- Network with one dominant interferer



Interference channel

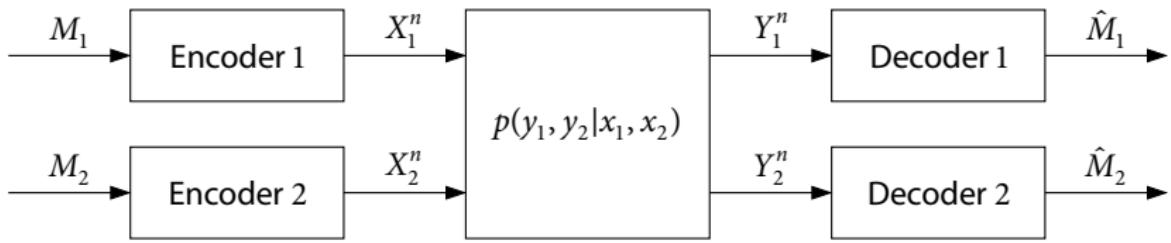


- **Capacity region:** Optimal tradeoff between the rates

$$R_1 = \frac{1}{n} \log |\text{supp}(M_1)|$$

$$R_2 = \frac{1}{n} \log |\text{supp}(M_2)|$$

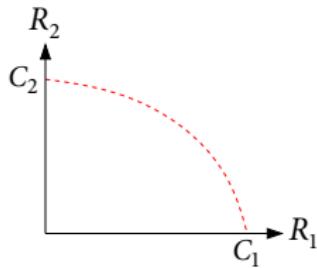
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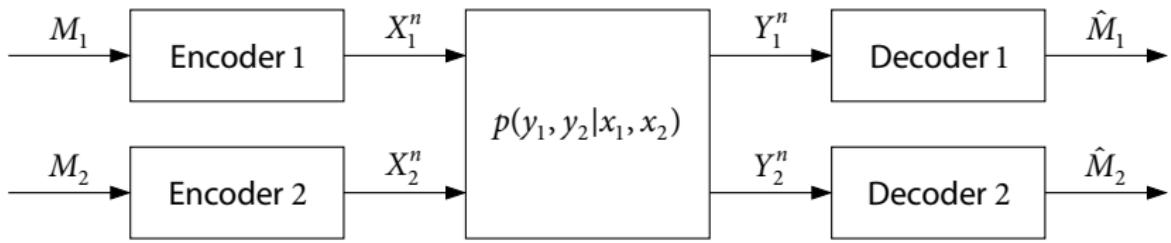
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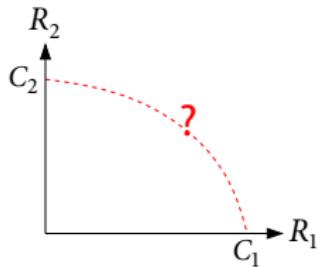
Interference channel



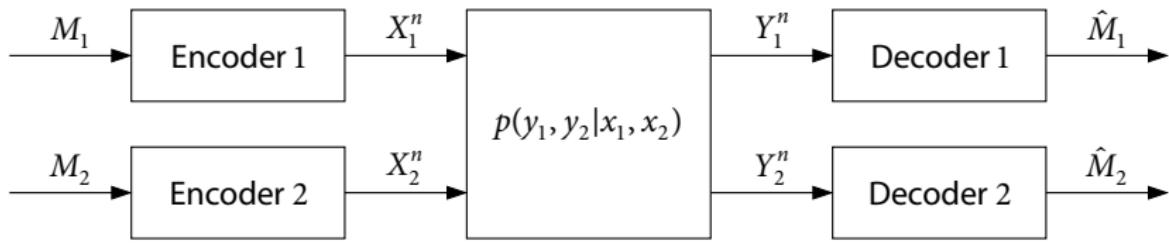
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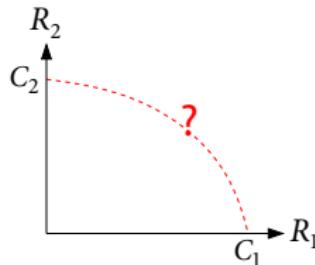
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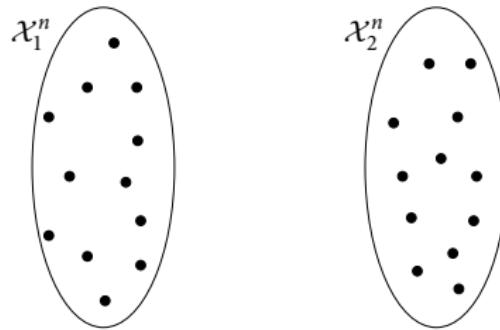
- Best inner bound: Han–Kobayashi (1981) and Nair–Xia–Yazdanpanah (2015)

Performance benchmark

- Highest rates achievable by point-to-point (P2P) random code ensembles

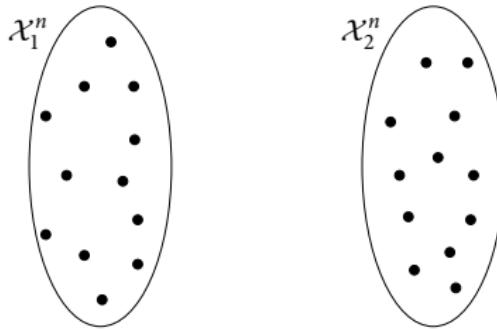
Performance benchmark

- Highest rates achievable by point-to-point (P2P) random code ensembles
- Random code ensemble
 - ▶ For each $m_1 \in [1 : 2^{nR_1}]$, generate $X_1^n(m_1) \sim \prod_{i=1}^n p_{X_1}(x_{1i})$
 - ▶ For each $m_2 \in [1 : 2^{nR_2}]$, generate $X_2^n(m_2) \sim \prod_{i=1}^n p_{X_2}(x_{2i})$



Performance benchmark

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- What is the optimal (MLD) tradeoff between achievable R_1 and R_2 ?

Why do we care?

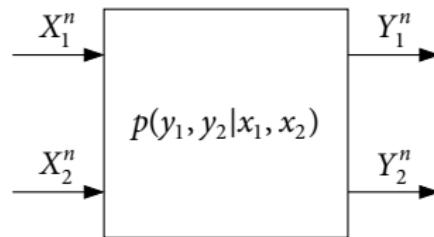
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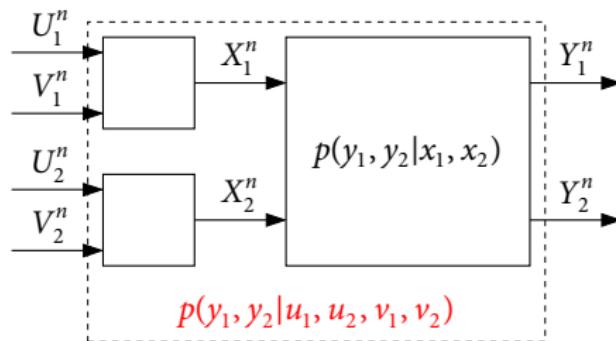
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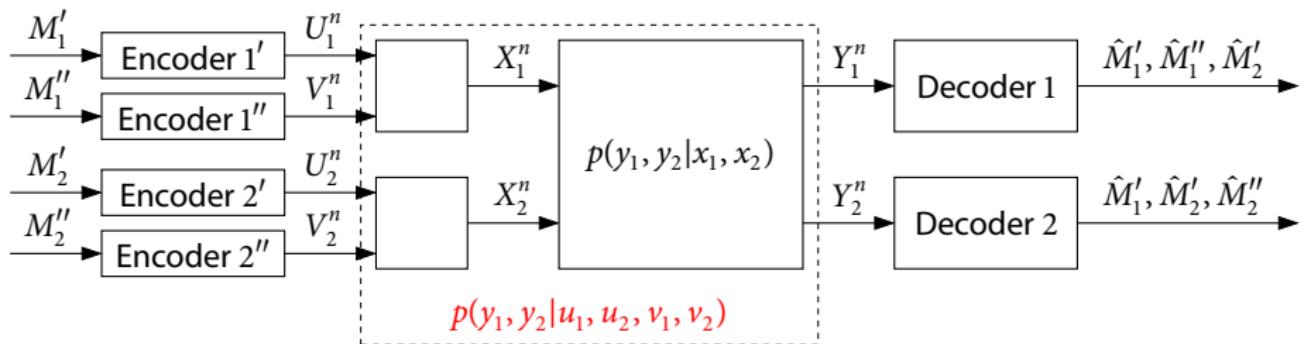
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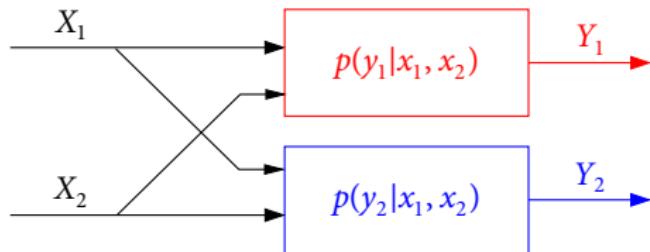


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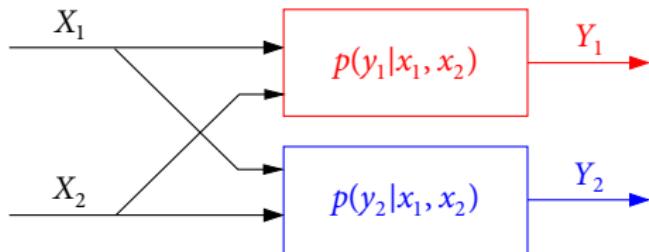
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Performance benchmark (Bandemer–El-Gamal–K 2012)



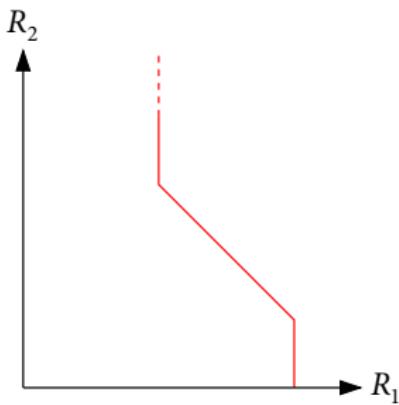
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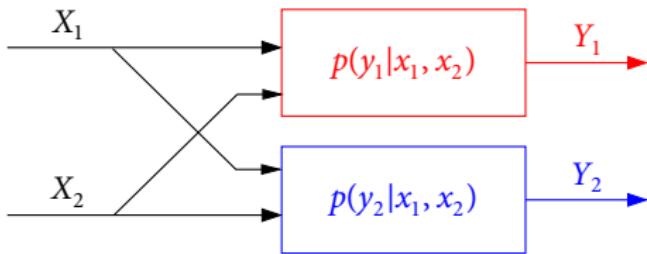
$$R_1 < I(X_1; Y_1 | X_2),$$
$$R_1 + R_2 < I(X_1, X_2; Y_1)$$

or

$$R_1 < I(X_1; Y_1)$$



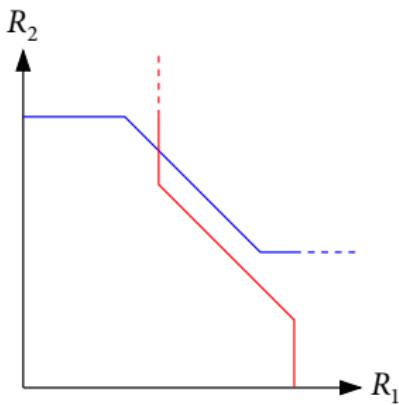
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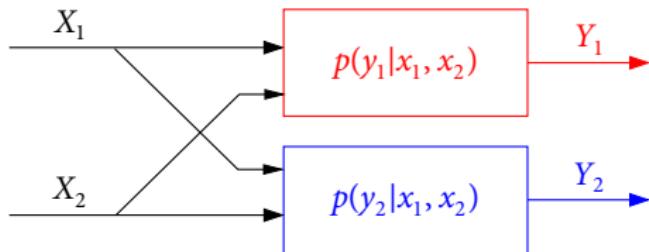
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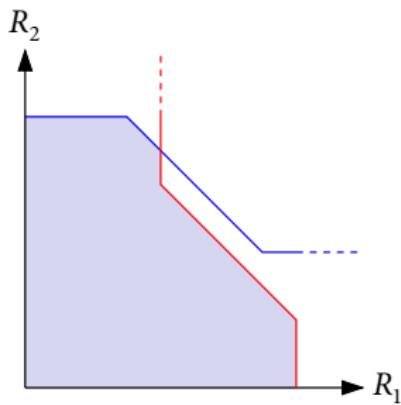
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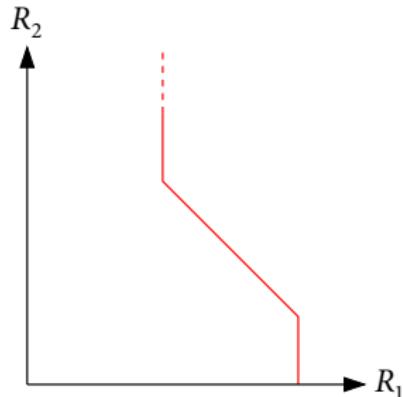


Maximum likelihood decoding \approx simultaneous decoding

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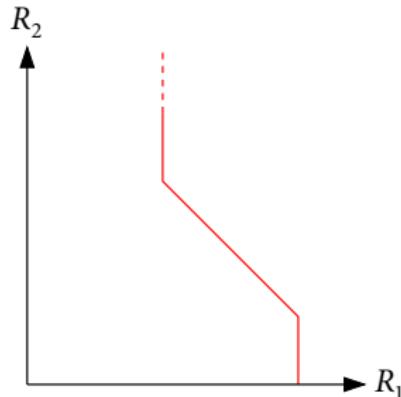


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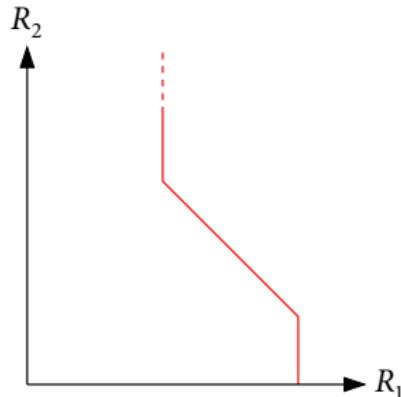
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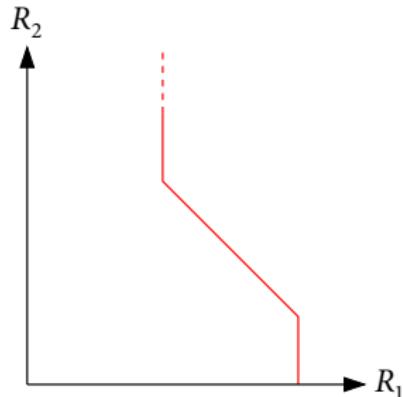
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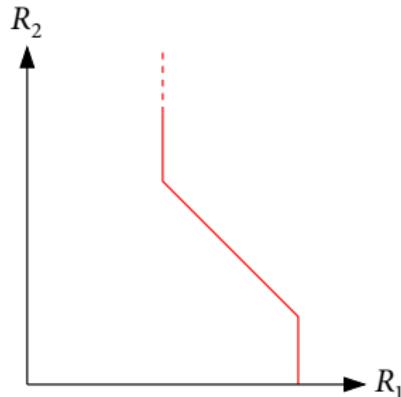
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Maximum likelihood decoding \approx simultaneous decoding

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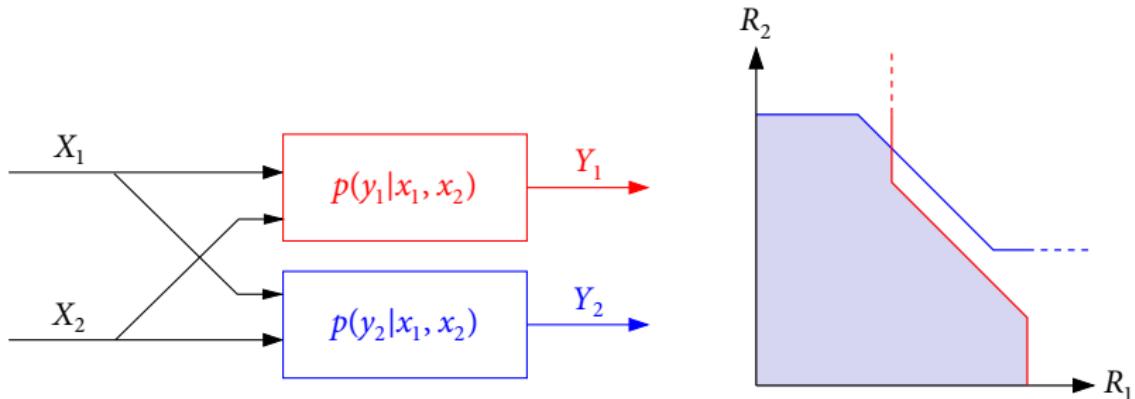
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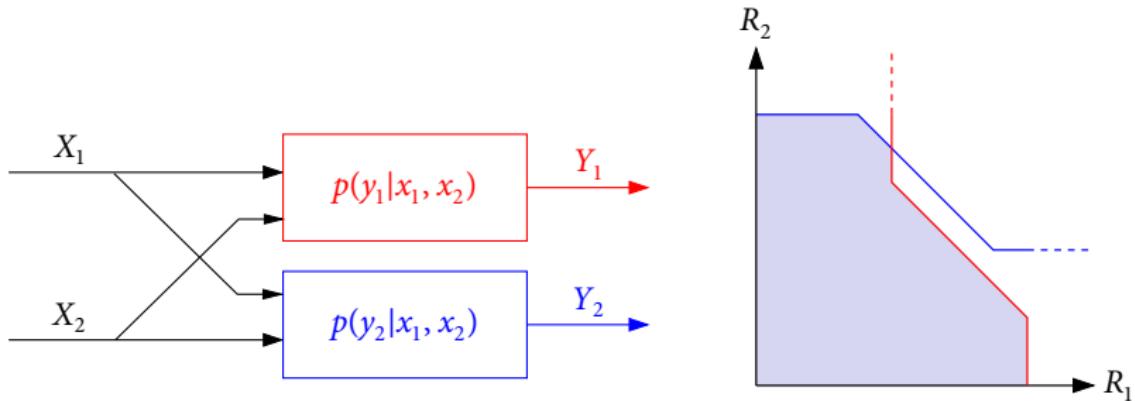


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- Multiuser detection: High complexity!

Low-complexity (implementable) alternatives

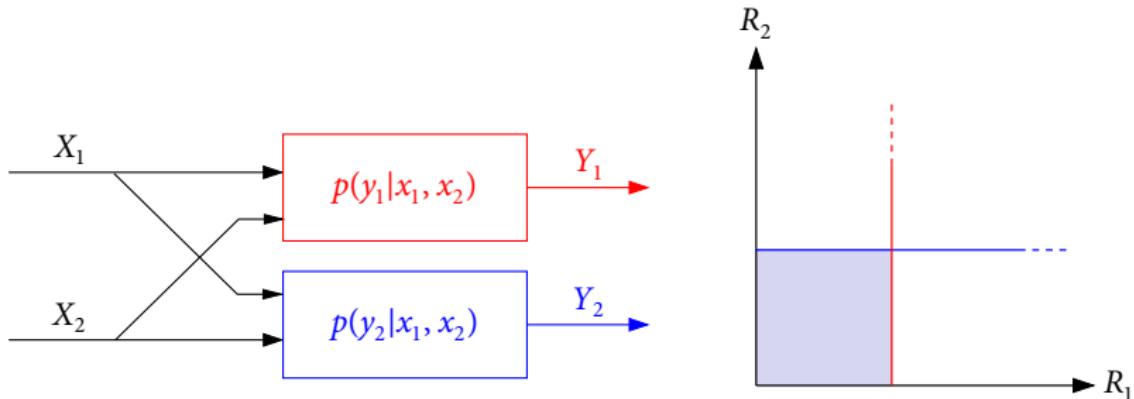


Low-complexity (implementable) alternatives



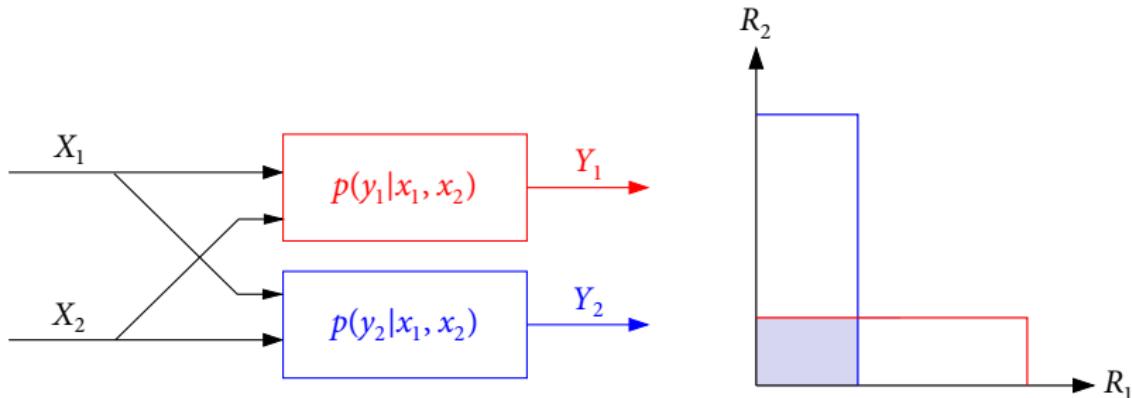
- P2P decoding

Low-complexity (implementable) alternatives



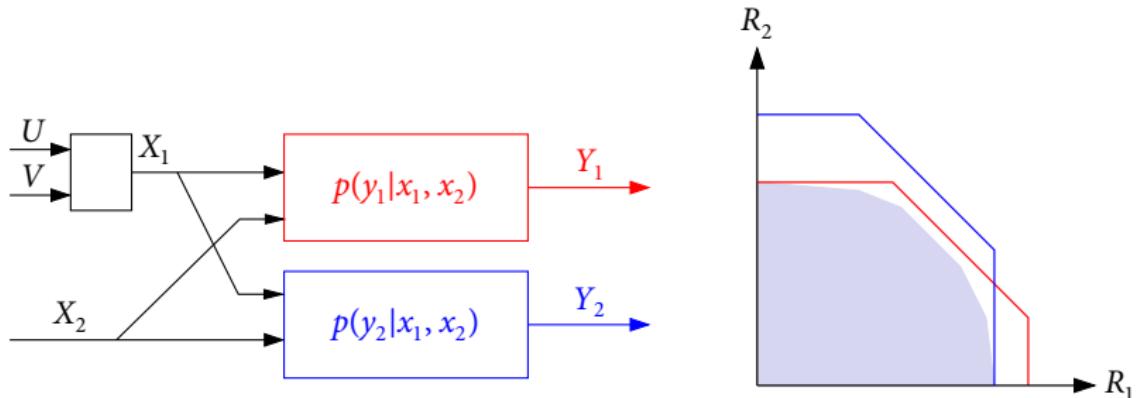
- P2P decoding
 - ▶ Treating interference as (Gaussian) noise: $R_1 < I(X_1; Y_1)$

Low-complexity (implementable) alternatives



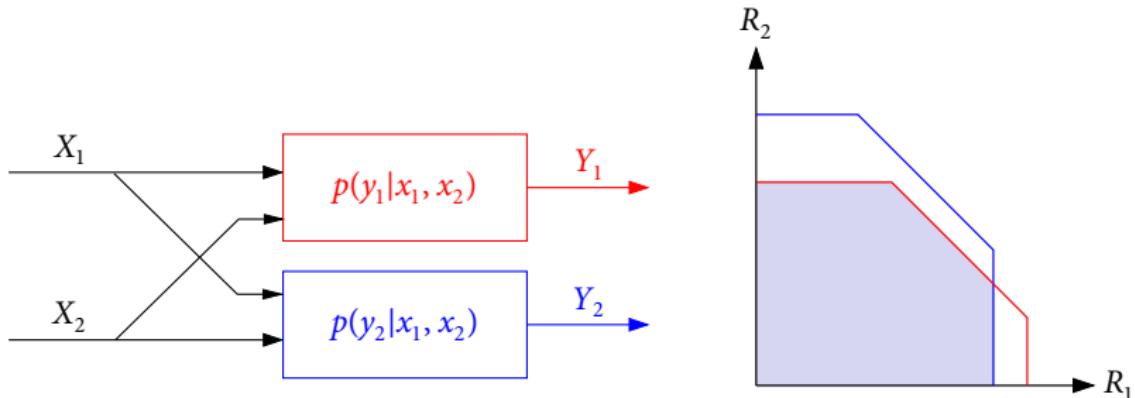
- P2P decoding
 - ▶ Treating interference as (Gaussian) noise: $R_1 < I(X_1; Y_1)$
 - ▶ Successive cancellation decoding: $R_2 < I(X_2; Y_1)$, $R_1 < I(X_1; Y_1|X_2)$

Low-complexity (implementable) alternatives



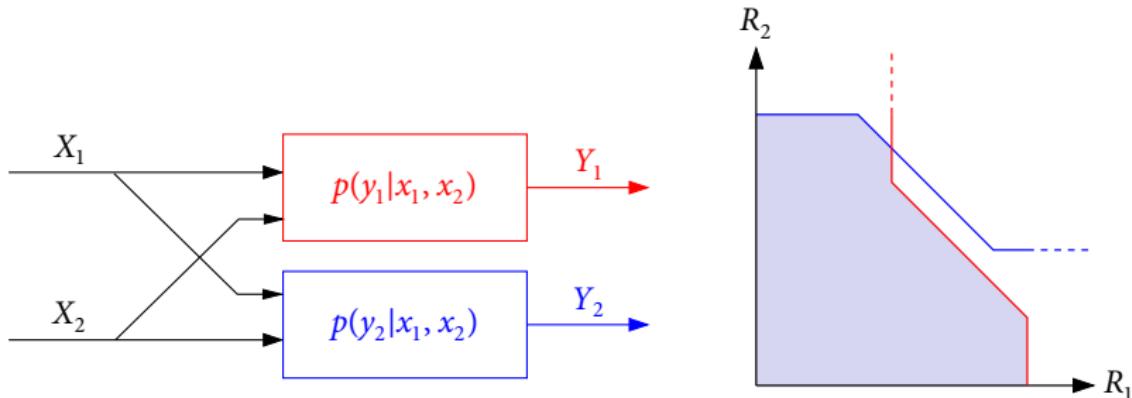
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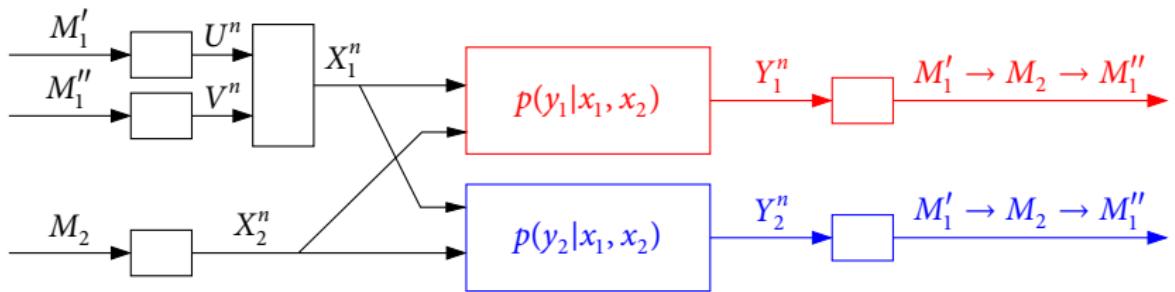
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 - ▶ Spatially coupled codes (Yedla, Nguyen, Pfister, and Narayanan 2011)
 - ▶ Polar codes (Wang and Şaşoğlu 2014)

Low-complexity (implementable) alternatives

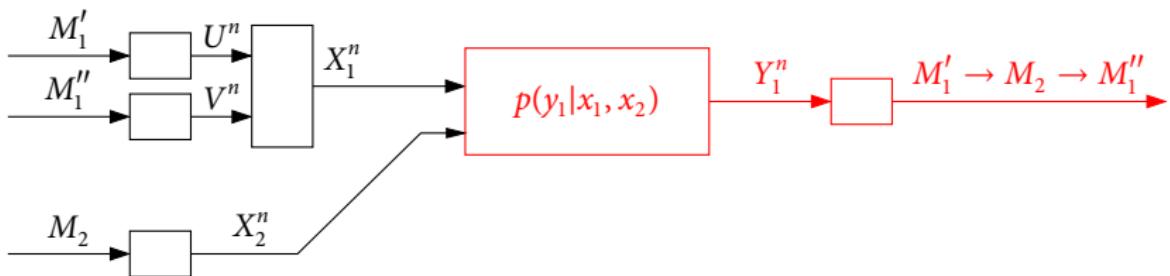


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A lesson from rate splitting (Grant et al. 2001)



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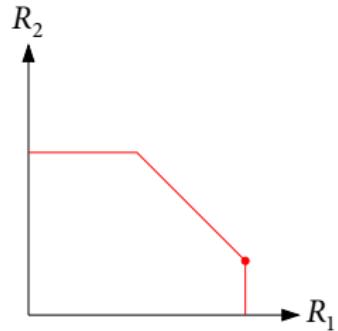


- Decoding at receiver 1:

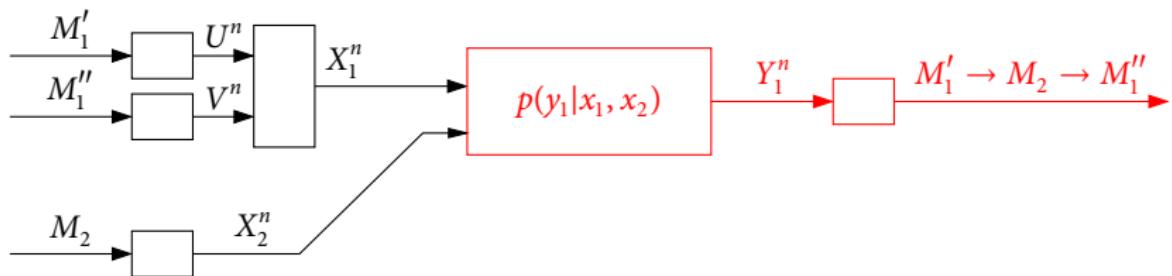
$$R'_1 < I(U; Y_1)$$

$$R_2 < I(X_2; Y_1 | U)$$

$$R''_1 < I(V; Y_1 | U, X_2)$$



A lesson from rate splitting (Grant et al. 2001)

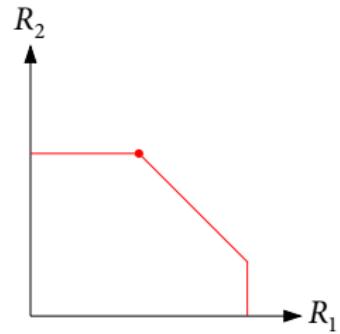


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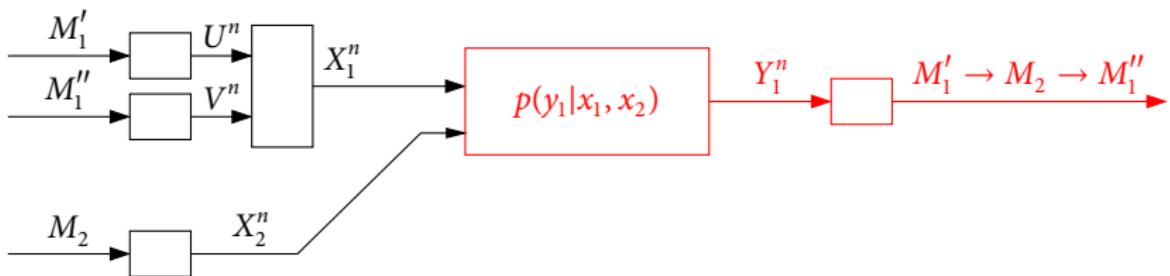
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$$R''_1 < I(V; Y_1 | U, X_2)$$



A lesson from rate splitting (Grant et al. 2001)

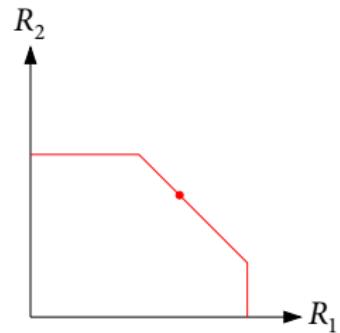


- Decoding at receiver 1:

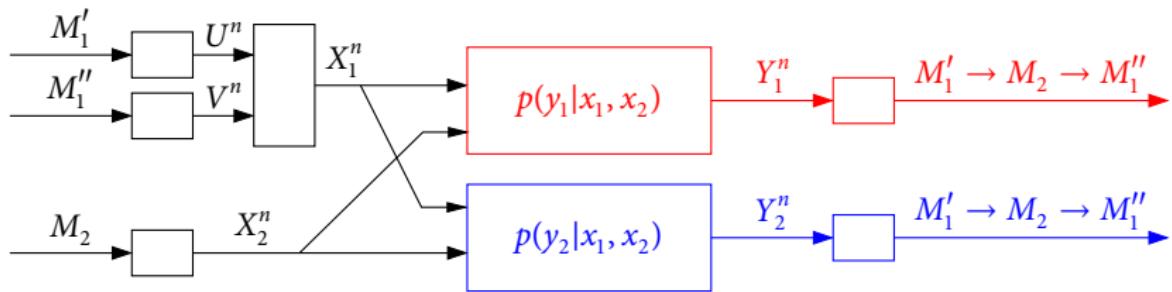
$$R'_1 < I(U; Y_1)$$

$$R_2 < I(X_2; Y_1 | U)$$

$$R''_1 < I(V; Y_1 | U, X_2)$$



A lesson from rate splitting (Grant et al. 2001)



- Decoding at receiver 1:

$$R'_1 < I(U; Y_1)$$

$$R_2 < I(X_2; Y_1 | U)$$

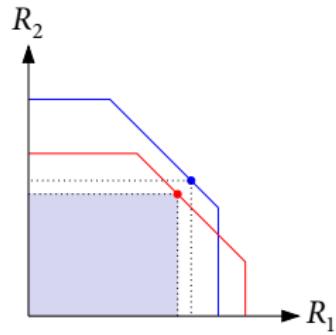
$$R''_1 < I(V; Y_1 | U, X_2)$$

- Decoding at receiver 2:

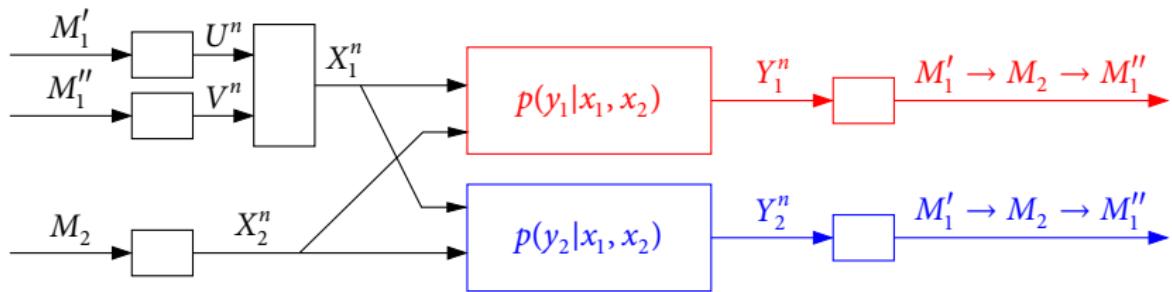
$$R'_1 < I(U; Y_2)$$

$$R_2 < I(X_2; Y_2 | U)$$

$$R''_1 < I(V; Y_2 | U, X_2)$$



A lesson from rate splitting (Grant et al. 2001)



- Decoding at receiver 1:

$$R'_1 < I(U; Y_1)$$

$$R_2 < I(X_2; Y_1 | U)$$

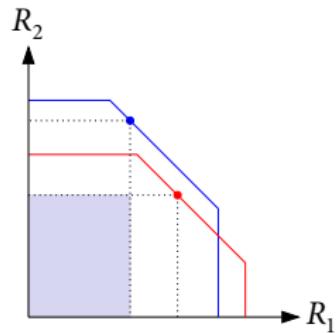
$$R''_1 < I(V; Y_1 | U, X_2)$$

- Decoding at receiver 2:

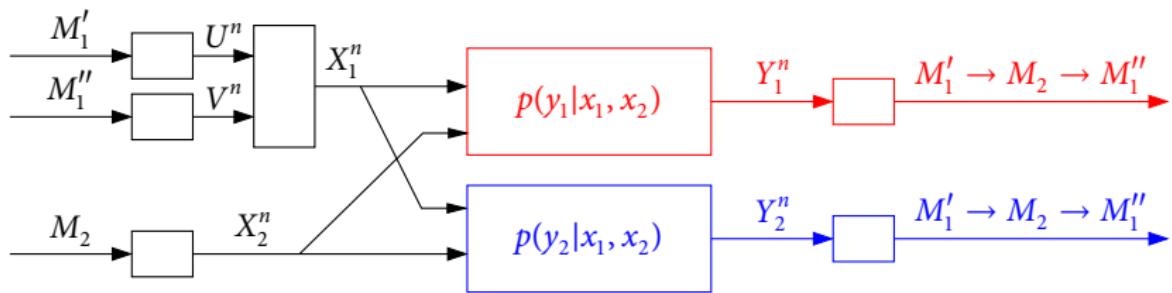
$$R'_1 < I(U; Y_2)$$

$$R_2 < I(X_2; Y_2 | U)$$

$$R''_1 < I(V; Y_2 | U, X_2)$$

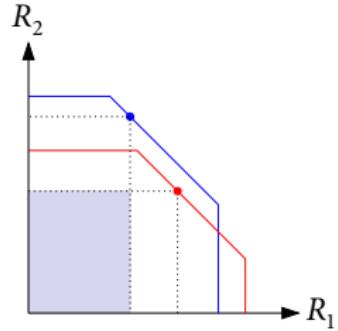


A lesson from rate splitting (Grant et al. 2001)

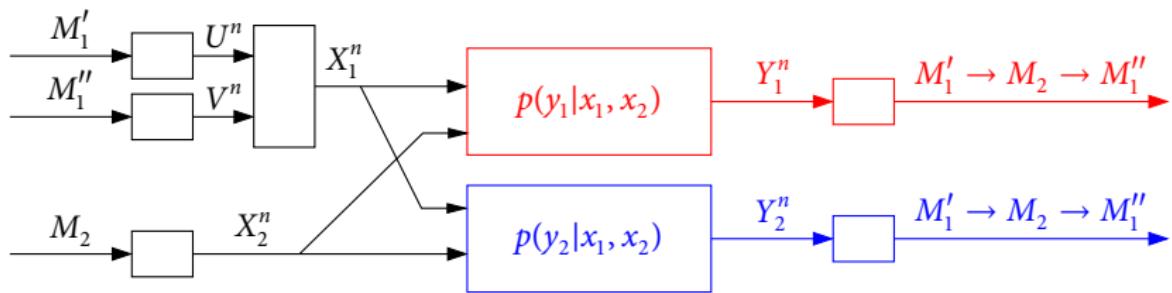


- Decoding at receiver 1: $R'_1 < I(U; Y_1)$, $R_2 < I(X_2; Y_1|U)$, $R''_1 < I(V; Y_1|U, X_2)$
- Decoding at receiver 2: $R'_1 < I(U; Y_2)$, $R_2 < I(X_2; Y_2|U)$, $R''_1 < I(V; Y_2|U, X_2)$
- Combined rate:

$$R_1 < \min_j I(U; Y_j) + \min_j I(V; Y_j|U, X_2)$$

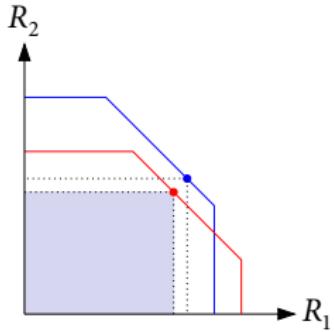


A lesson from rate splitting (Grant et al. 2001)

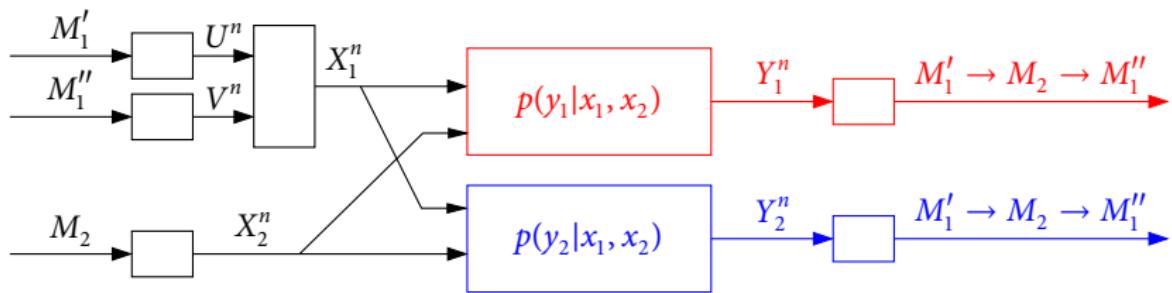


- Decoding at receiver 1: $R'_1 < I(U; Y_1)$, $R_2 < I(X_2; Y_1|U)$, $R''_1 < I(V; Y_1|U, X_2)$
- Decoding at receiver 2: $R'_1 < I(U; Y_2)$, $R_2 < I(X_2; Y_2|U)$, $R''_1 < I(V; Y_2|U, X_2)$
- Combined rate:

$$\begin{aligned} R_1 &< \min_j I(U; Y_j) + \min_j I(V; Y_j|U, X_2) \\ &< \min_j [I(U; Y_j) + I(V; Y_j|U, X_2)] \end{aligned}$$



A lesson from rate splitting (Grant et al. 2001)

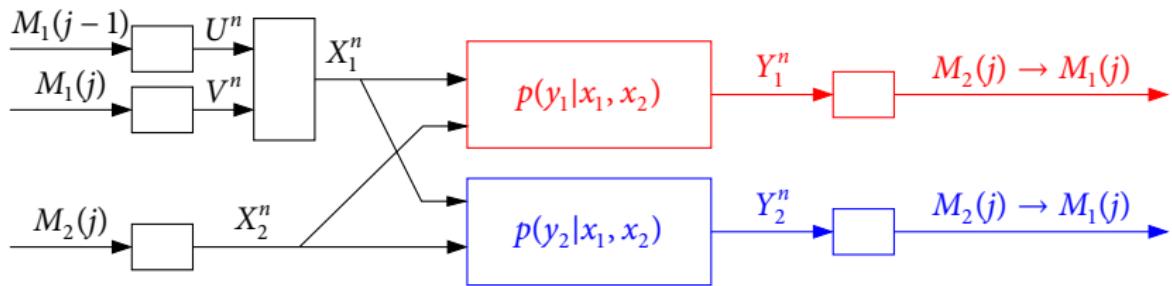


- Decoding at receiver 1: $R'_1 < I(U; Y_1)$, $R_2 < I(X_2; Y_1|U)$, $R''_1 < I(V; Y_1|U, X_2)$
- Decoding at receiver 2: $R'_1 < I(U; Y_2)$, $R_2 < I(X_2; Y_2|U)$, $R''_1 < I(V; Y_2|U, X_2)$
- Combined rate:

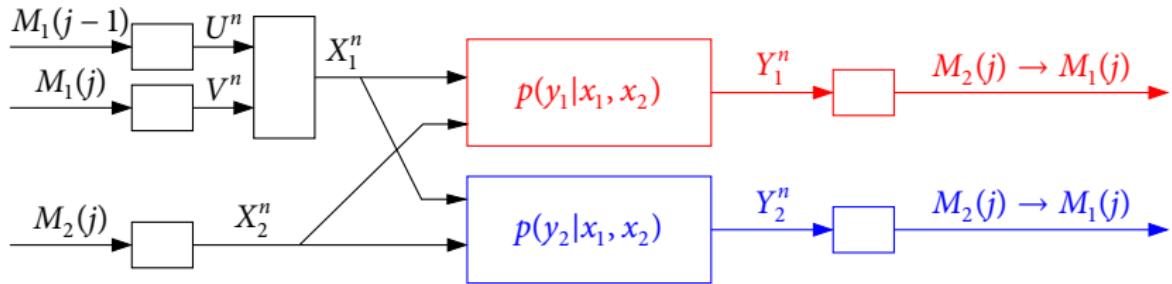
$$\begin{aligned} R_1 &< \min_j I(U; Y_j) + \min_j I(V; Y_j|U, X_2) \\ &< \min_j [I(U; Y_j) + I(V; Y_j|U, X_2)] \end{aligned}$$

- Key to achieving the SD performance: Switch the order of sum and min!

Sliding-window superposition coding (Wang et al. 2014)

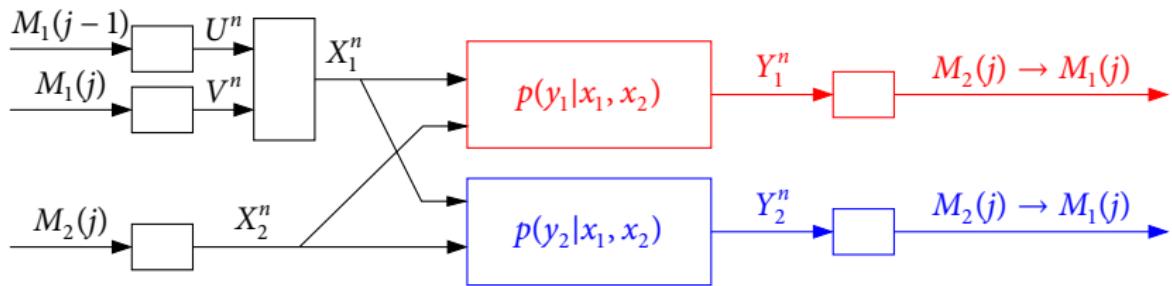


Sliding-window superposition coding (Wang et al. 2014)



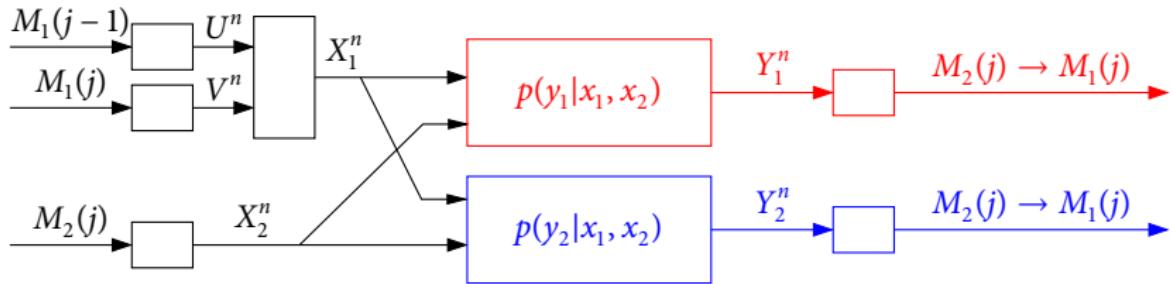
| Block | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|-------|----------|----------|----------|----------|----------|----------|----------|
| U | | | | | | | |
| V | | | | | | | |
| X_2 | $M_2(1)$ | $M_2(2)$ | $M_2(3)$ | $M_2(4)$ | $M_2(5)$ | $M_2(6)$ | $M_2(7)$ |

Sliding-window superposition coding (Wang et al. 2014)



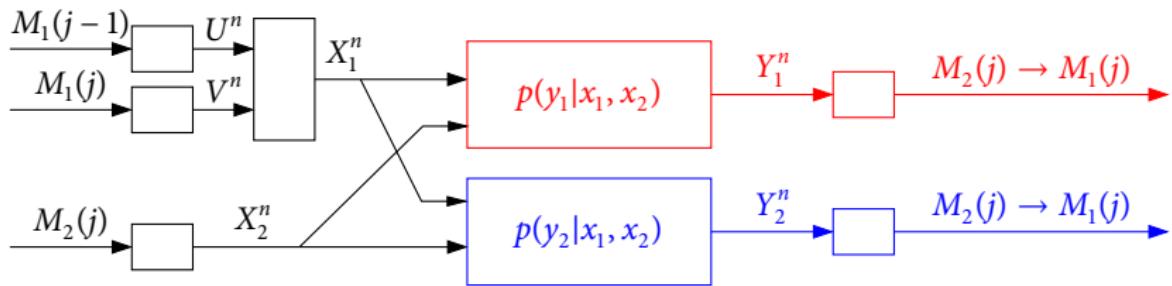
| Block | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|-------|----------|----------|----------|----------|----------|----------|----------|
| U | | $M_1(1)$ | | | | | |
| V | $M_1(1)$ | | | | | | |
| X_2 | $M_2(1)$ | $M_2(2)$ | $M_2(3)$ | $M_2(4)$ | $M_2(5)$ | $M_2(6)$ | $M_2(7)$ |

Sliding-window superposition coding (Wang et al. 2014)



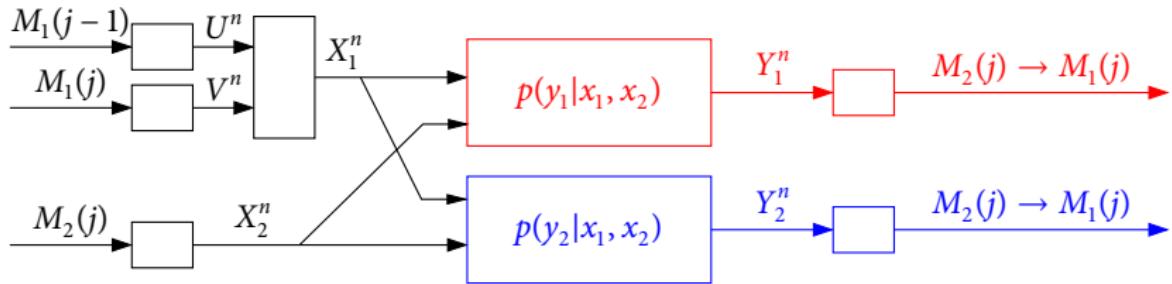
| Block | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|-------|----------|----------|----------|----------|----------|----------|----------|
| U | | $M_1(1)$ | $M_1(2)$ | | | | |
| V | $M_1(1)$ | $M_1(2)$ | | | | | |
| X_2 | $M_2(1)$ | $M_2(2)$ | $M_2(3)$ | $M_2(4)$ | $M_2(5)$ | $M_2(6)$ | $M_2(7)$ |

Sliding-window superposition coding (Wang et al. 2014)



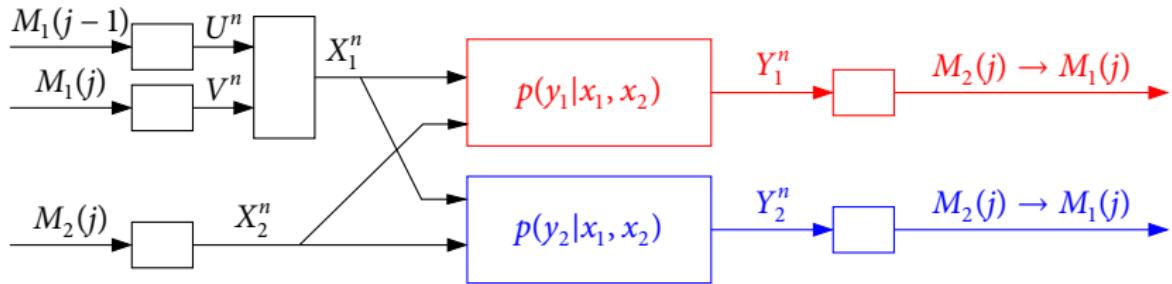
| Block | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|-------|----------|----------|----------|----------|----------|----------|----------|
| U | | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | | | |
| V | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | | | | |
| X_2 | $M_2(1)$ | $M_2(2)$ | $M_2(3)$ | $M_2(4)$ | $M_2(5)$ | $M_2(6)$ | $M_2(7)$ |

Sliding-window superposition coding (Wang et al. 2014)



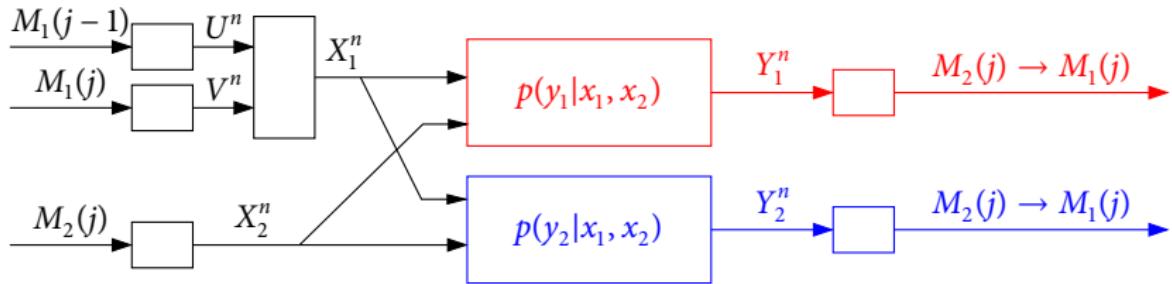
| Block | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|-------|----------|----------|----------|----------|----------|----------|----------|
| U | | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | | |
| V | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | | | |
| X_2 | $M_2(1)$ | $M_2(2)$ | $M_2(3)$ | $M_2(4)$ | $M_2(5)$ | $M_2(6)$ | $M_2(7)$ |

Sliding-window superposition coding (Wang et al. 2014)



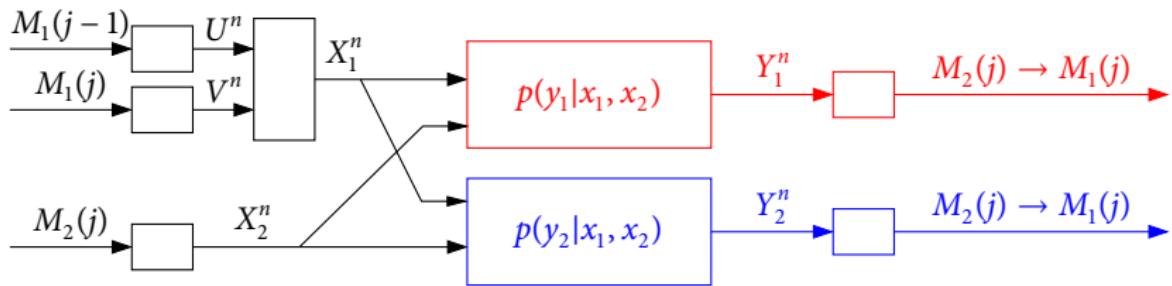
| Block | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|-------|----------|----------|----------|----------|----------|----------|----------|
| U | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | | |
| V | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | | |
| X_2 | $M_2(1)$ | $M_2(2)$ | $M_2(3)$ | $M_2(4)$ | $M_2(5)$ | $M_2(6)$ | $M_2(7)$ |

Sliding-window superposition coding (Wang et al. 2014)



| Block | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|-------|----------|----------|----------|----------|----------|----------|----------|
| U | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ | |
| V | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ | |
| X_2 | $M_2(1)$ | $M_2(2)$ | $M_2(3)$ | $M_2(4)$ | $M_2(5)$ | $M_2(6)$ | $M_2(7)$ |

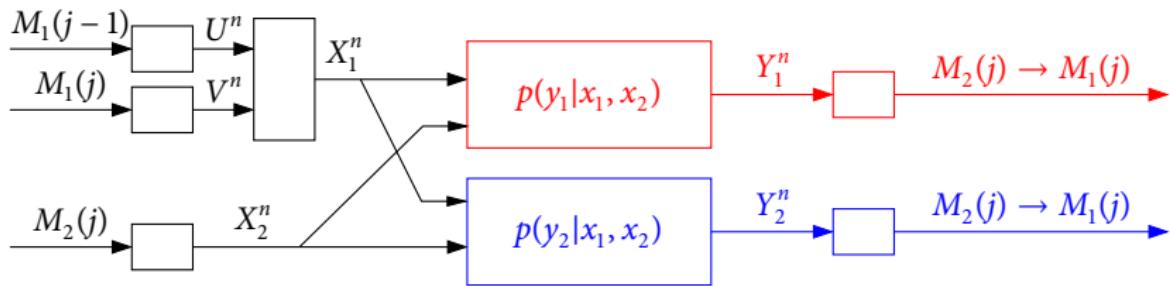
Sliding-window superposition coding (Wang et al. 2014)



| Block | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|-------|----------|----------|----------|----------|----------|----------|----------|
| U | | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ |
| V | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ | |
| X_2 | $M_2(1)$ | $M_2(2)$ | $M_2(3)$ | $M_2(4)$ | $M_2(5)$ | $M_2(6)$ | $M_2(7)$ |

- **Block Markov coding:** As in relaying and feedback communication

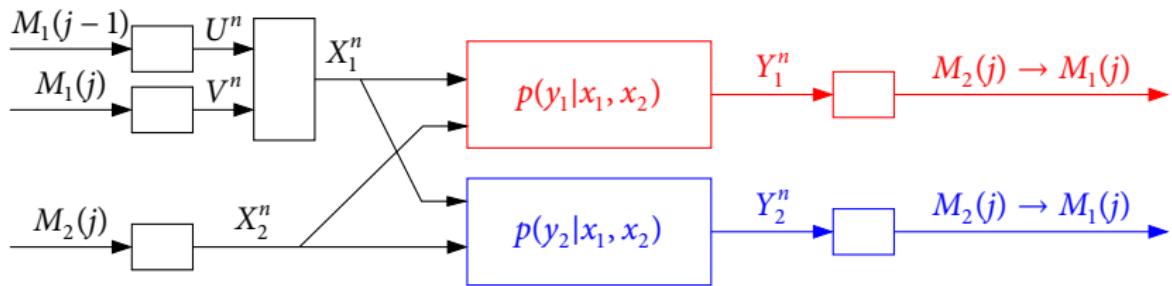
Sliding-window superposition coding (Wang et al. 2014)



| Block | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|-------|----------|----------|----------|----------|----------|----------|----------|
| U | | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ |
| V | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ | |
| X_2 | $M_2(1)$ | $M_2(2)$ | $M_2(3)$ | $M_2(4)$ | $M_2(5)$ | $M_2(6)$ | $M_2(7)$ |

- **Block Markov coding:** As in relaying and feedback communication
- **Superposition coding:** But without rate splitting

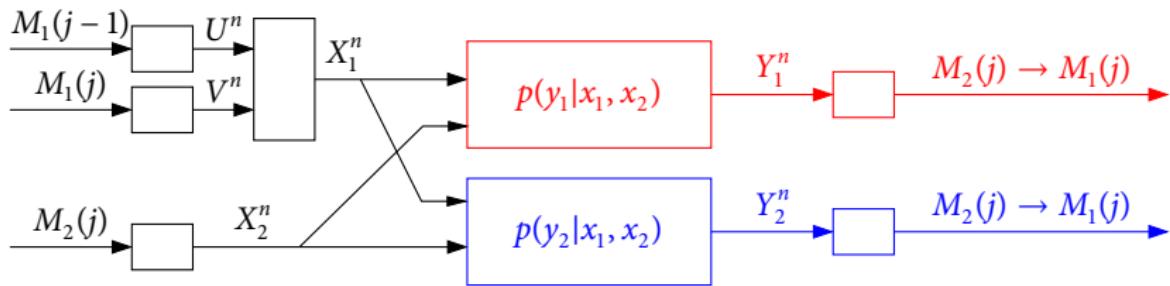
Sliding-window superposition coding (Wang et al. 2014)



| Block | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|-------|----------|----------|----------|----------|----------|----------|----------|
| U | | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ |
| V | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ | |
| X_2 | $M_2(1)$ | $M_2(2)$ | $M_2(3)$ | $M_2(4)$ | $M_2(5)$ | $M_2(6)$ | $M_2(7)$ |

- **Block Markov coding:** As in relaying and feedback communication
- **Superposition coding:** But without rate splitting
- **Staggered (asynchronous) transmission:** cf. EV-DO rev A, D-BLAST

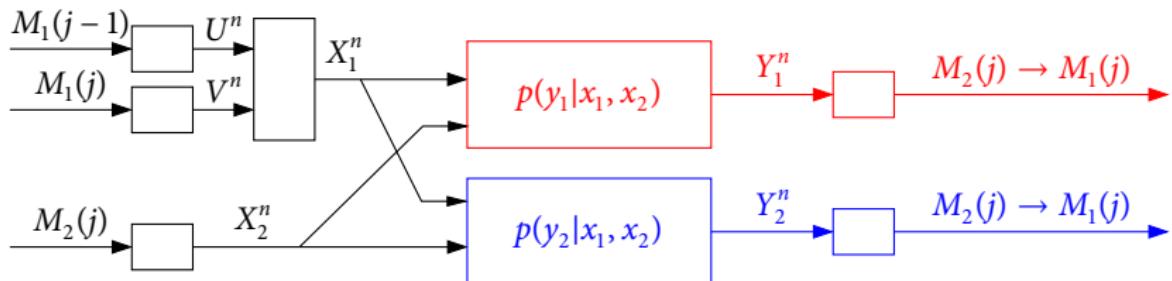
Sliding-window superposition coding (Wang et al. 2014)



| Block | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|-------|----------|----------|----------|----------|----------|----------|----------|
| U | | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ |
| V | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ | |
| X_2 | $M_2(1)$ | $M_2(2)$ | $M_2(3)$ | $M_2(4)$ | $M_2(5)$ | $M_2(6)$ | $M_2(7)$ |

- Sliding-window decoding

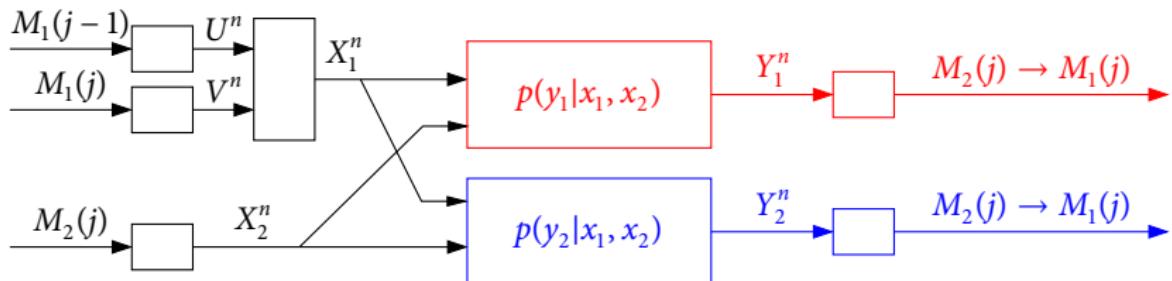
Sliding-window superposition coding (Wang et al. 2014)



| Block | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|-------|---|----------|----------|----------|----------|----------|----------|
| U | | | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ |
| V | | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ | |
| X_2 | | $M_2(2)$ | $M_2(3)$ | $M_2(4)$ | $M_2(5)$ | $M_2(6)$ | $M_2(7)$ |

- Sliding-window decoding
- Successive cancellation decoding

Sliding-window superposition coding (Wang et al. 2014)

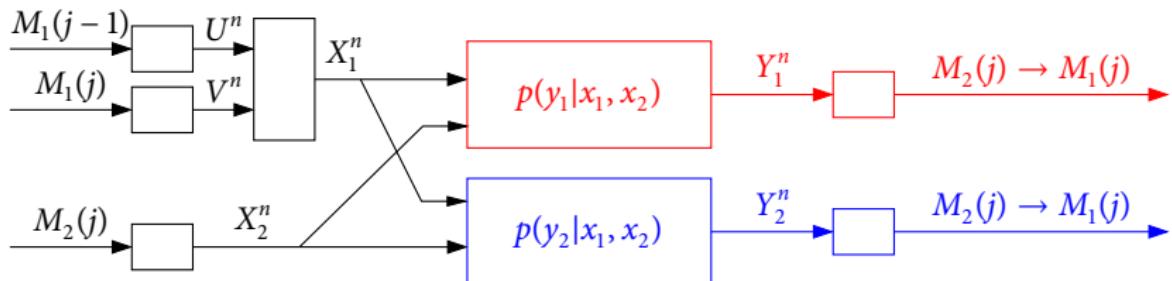


| Block | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|-------|---|---|----------|----------|----------|----------|----------|
| U | | | | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ |
| V | | | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ | |
| X_2 | | | $M_2(3)$ | $M_2(4)$ | $M_2(5)$ | $M_2(6)$ | $M_2(7)$ |

- Sliding-window decoding
- Successive cancellation decoding

$$R_2 < I(X_2; Y_j | U),$$

Sliding-window superposition coding (Wang et al. 2014)



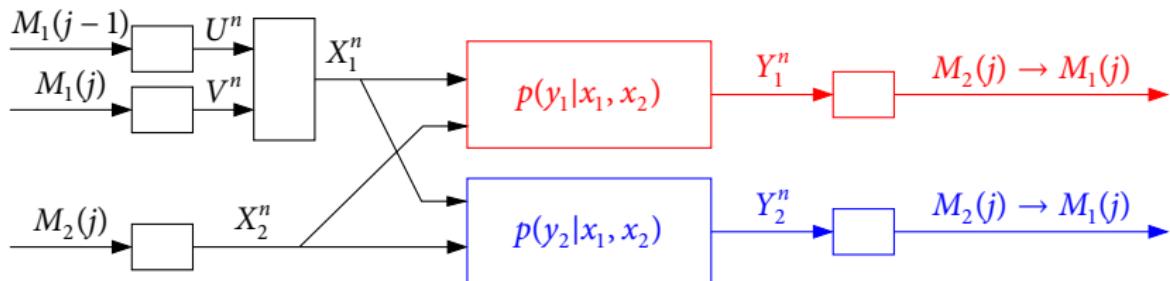
| Block | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|-------|---|---|----------|----------|----------|----------|----------|
| U | | | | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ |
| V | | | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ | |
| X_2 | | | | $M_2(4)$ | $M_2(5)$ | $M_2(6)$ | $M_2(7)$ |

- Sliding-window decoding
- Successive cancellation decoding

$$R_2 < I(X_2; Y_j | U),$$

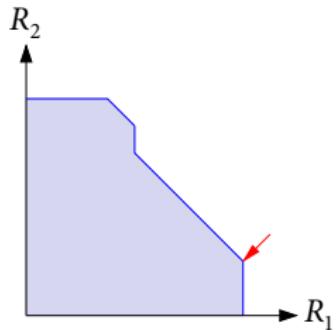
$$R_1 < I(U; Y_j) + I(V; Y_j | U, X_2)$$

Sliding-window superposition coding (Wang et al. 2014)

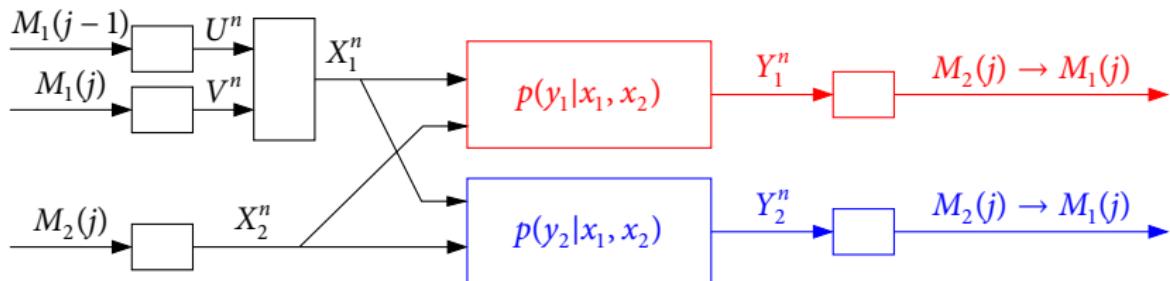


| Block | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|-------|----------|----------|----------|----------|----------|----------|----------|
| U | | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ |
| V | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ | |
| X_2 | $M_2(1)$ | $M_2(2)$ | $M_2(3)$ | $M_2(4)$ | $M_2(5)$ | $M_2(6)$ | $M_2(7)$ |

- Every corner point: different decoding orders

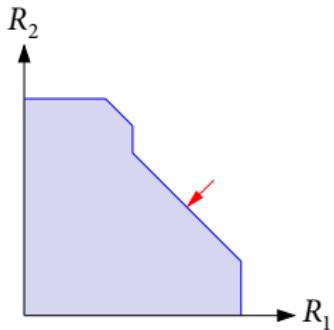


Sliding-window superposition coding (Wang et al. 2014)

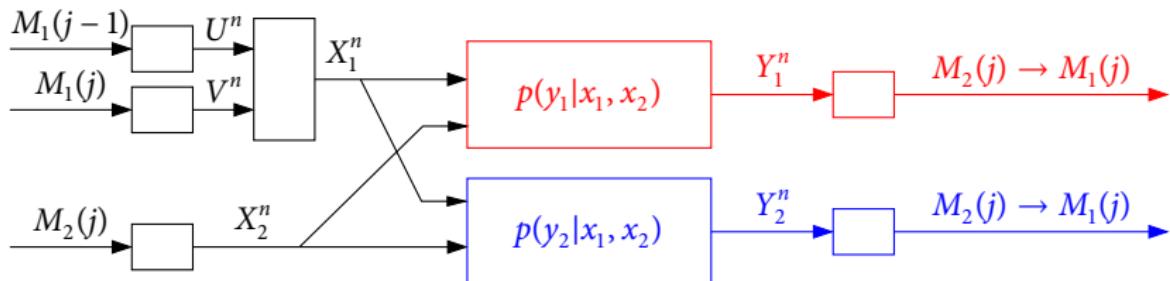


| Block | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|-------|----------|----------|----------|----------|----------|----------|----------|
| U | | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ |
| V | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ | |
| X_2 | $M_2(1)$ | $M_2(2)$ | $M_2(3)$ | $M_2(4)$ | $M_2(5)$ | $M_2(6)$ | $M_2(7)$ |

- Every corner point: different decoding orders
- Every point: time sharing or more superposition layers

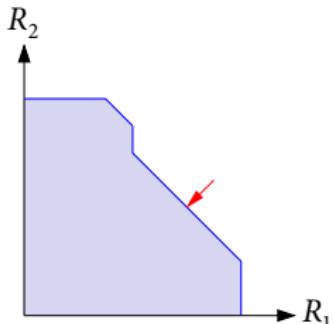


Sliding-window superposition coding (Wang et al. 2014)

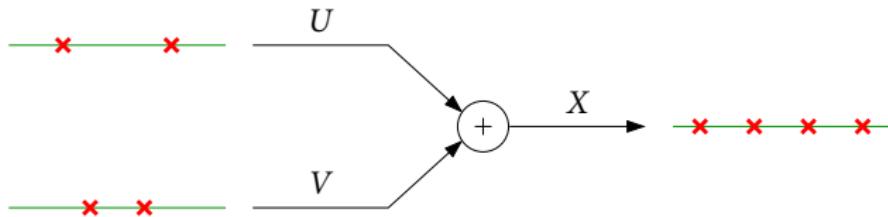


| Block | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|---------|------------|------------|------------|------------|------------|------------|------------|
| \$U\$ | | \$M_1(1)\$ | \$M_1(2)\$ | \$M_1(3)\$ | \$M_1(4)\$ | \$M_1(5)\$ | \$M_1(6)\$ |
| \$V\$ | \$M_1(1)\$ | \$M_1(2)\$ | \$M_1(3)\$ | \$M_1(4)\$ | \$M_1(5)\$ | \$M_1(6)\$ | |
| \$X_2\$ | \$M_2(1)\$ | \$M_2(2)\$ | \$M_2(3)\$ | \$M_2(4)\$ | \$M_2(5)\$ | \$M_2(6)\$ | \$M_2(7)\$ |

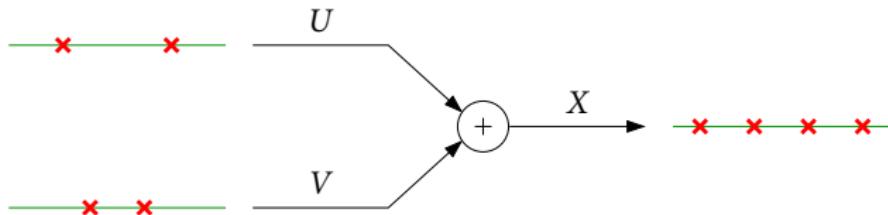
- Every corner point: different decoding orders
- Every point: time sharing or more superposition layers
- General theory for arbitrary number of users (Wang 2015)



Towards coded modulation



Towards coded modulation



| | |
|-----|-------|
| U | M' |
| V | M'' |

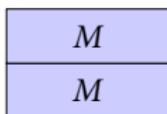
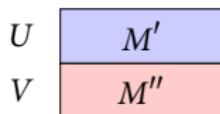
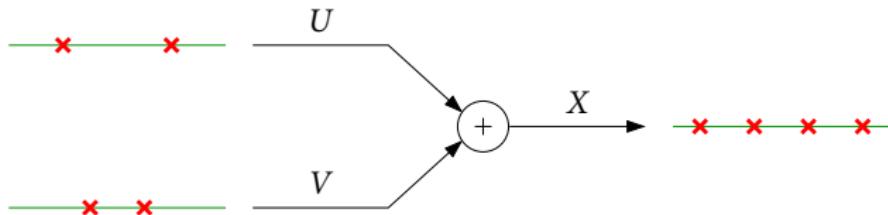
Multilevel coding (MLC)

$$R' < I(U; Y)$$

$$R'' < I(V; Y|U)$$

Short, nonuniversal

Towards coded modulation



Multilevel coding (MLC)

Bit-interleaved coded
modulation (BICM)

$$R' < I(U; Y)$$

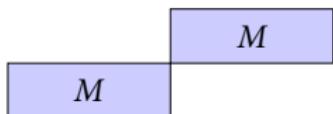
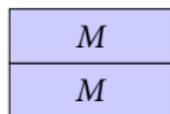
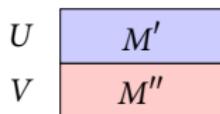
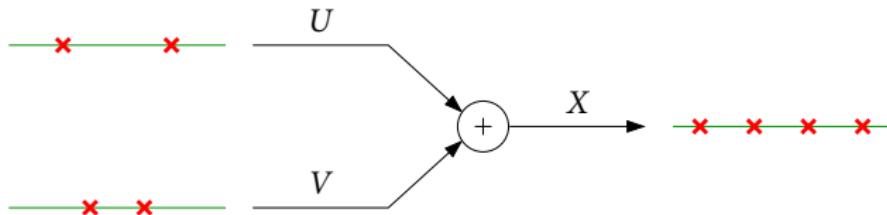
$$R < I(U; Y) + I(V; Y)$$

$$R'' < I(V; Y|U)$$

Short, nonuniversal

Other layers as noise

Towards coded modulation



Multilevel coding (MLC)

$$R' < I(U; Y)$$

$$R'' < I(V; Y|U)$$

Short, nonuniversal

Bit-interleaved coded modulation (BICM)

$$R < I(U; Y) + I(V; Y)$$

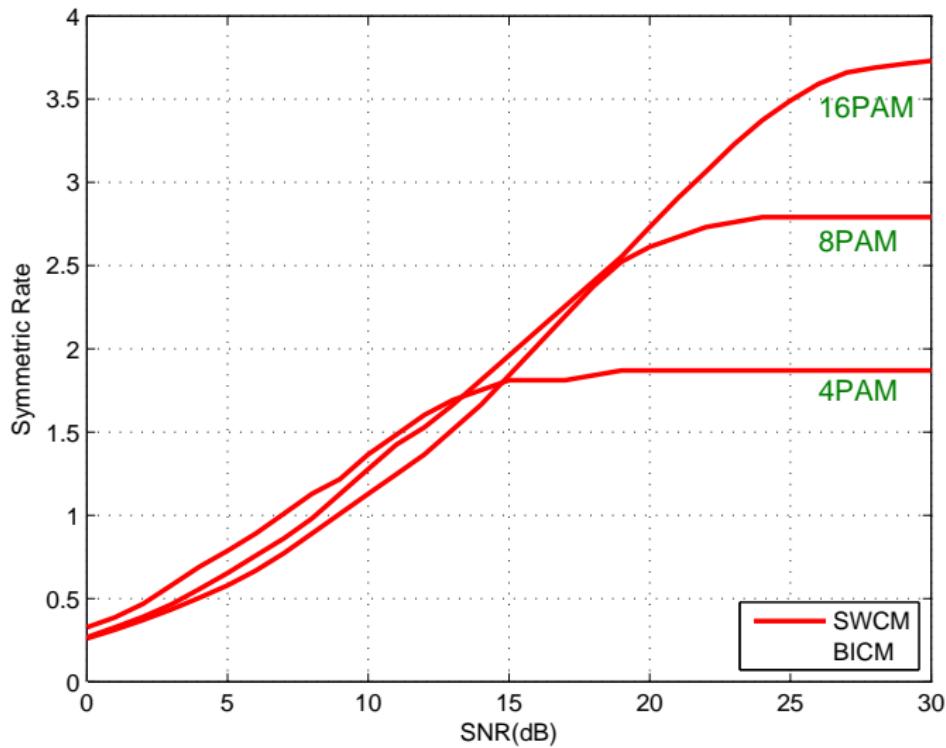
Other layers as noise

Sliding-window coded modulation (SWCM)

$$\begin{aligned} R &< I(U; Y) + I(V; Y|U) \\ &= I(X; Y) \end{aligned}$$

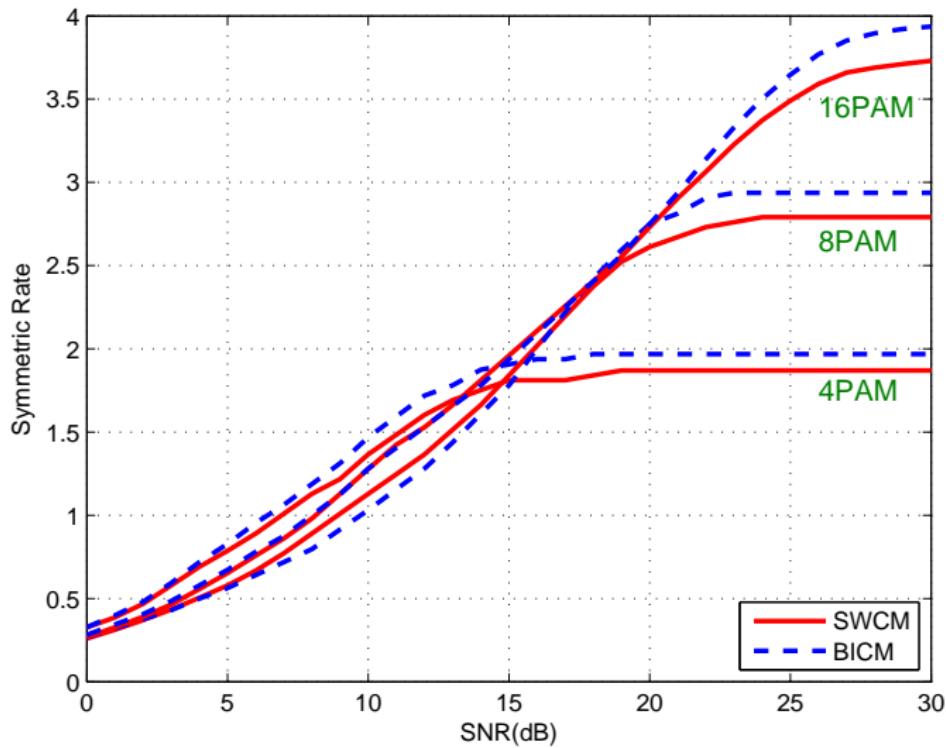
Error prop., rate loss

P2P performance



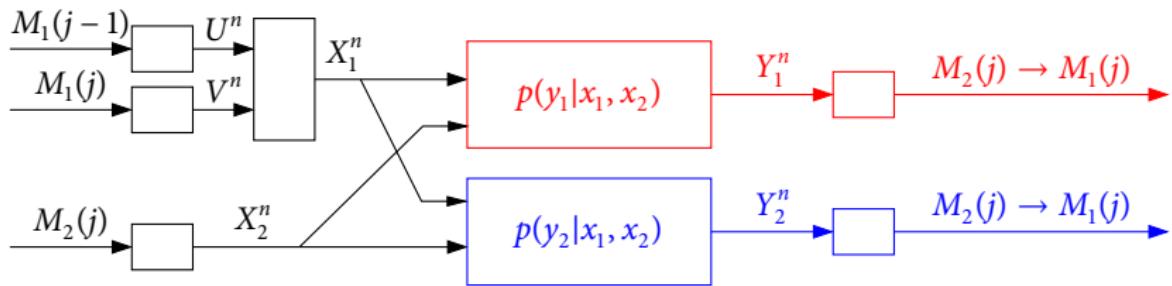
LTE turbo code / ≤ 8 -iteration LOG-MAP decoding at $b = 20$, $n = 2048$, BLER = 0.1

P2P performance



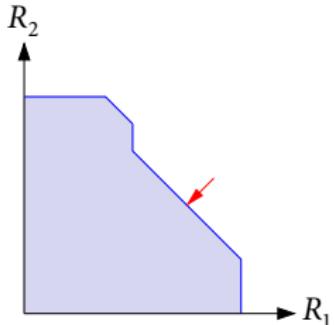
LTE turbo code / ≤ 8 -iteration LOG-MAP decoding at $b = 20$, $n = 2048$, BLER = 0.1

Back to interference mitigation

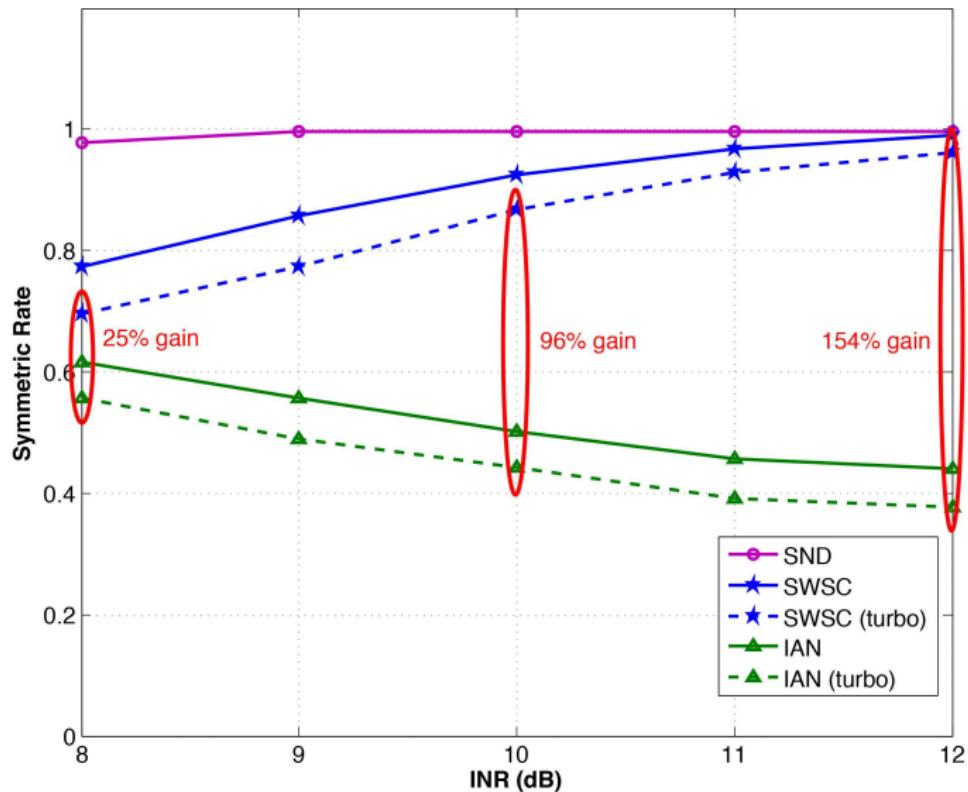


| Block | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|-------|----------|----------|----------|----------|----------|----------|----------|
| U | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ | |
| V | $M_1(1)$ | $M_1(2)$ | $M_1(3)$ | $M_1(4)$ | $M_1(5)$ | $M_1(6)$ | |
| X_2 | $M_2(1)$ | $M_2(2)$ | $M_2(3)$ | $M_2(4)$ | $M_2(5)$ | $M_2(6)$ | $M_2(7)$ |

- Every corner point: different decoding orders
- Every point: time sharing or more superposition layers
- General theory for arbitrary number of users (Wang 2015)

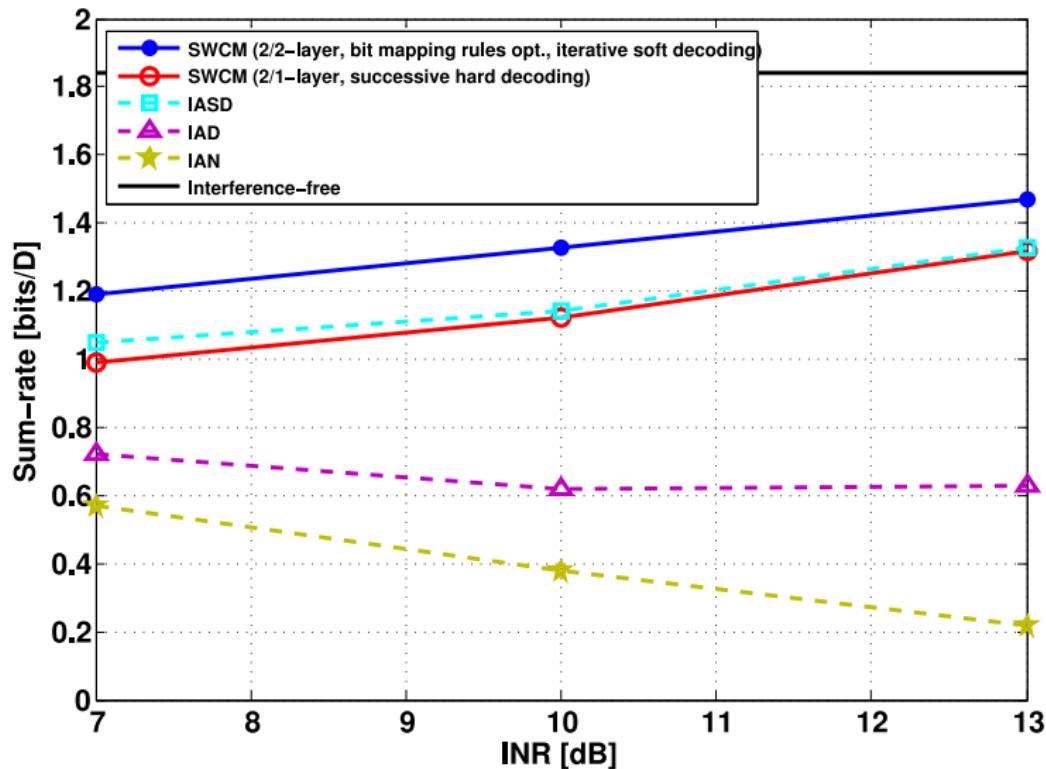


Gaussian channel performance (Park–K–Wang 2014)



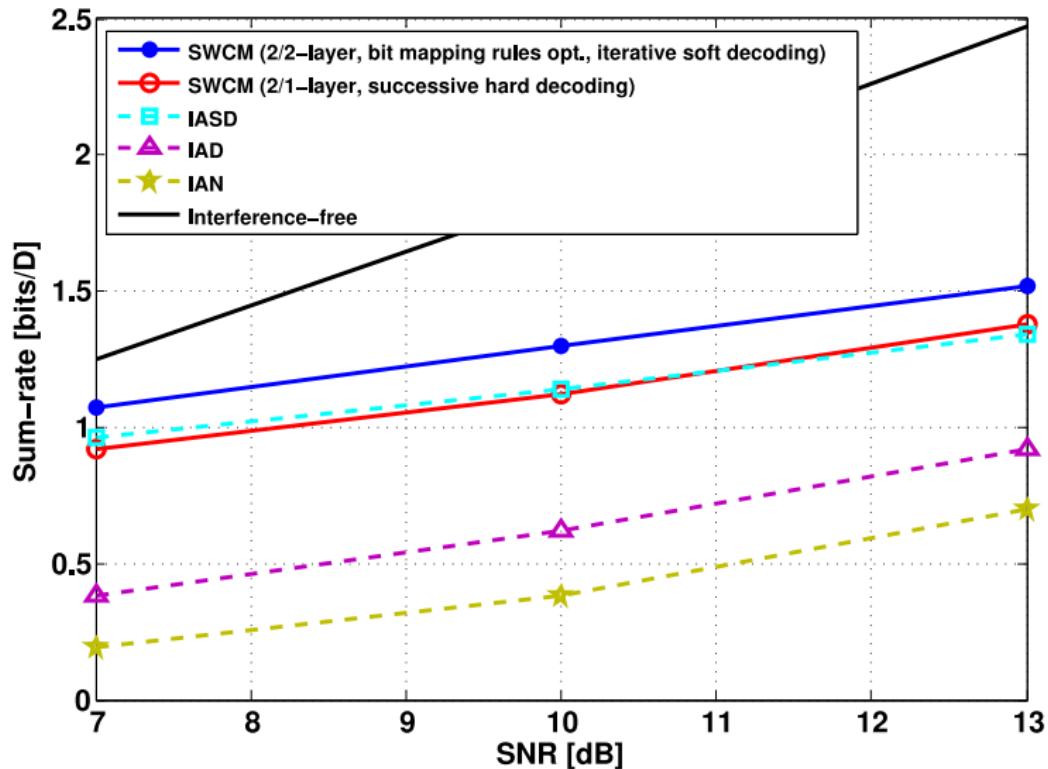
LTE turbo code / ≤ 8 -iteration LOG-MAP decoding at $b = 20$, $n = 2048$, BLER = 0.1, SNR = 10 dB

OFDM/Ped-B channel performance (Kim et al. 2015)



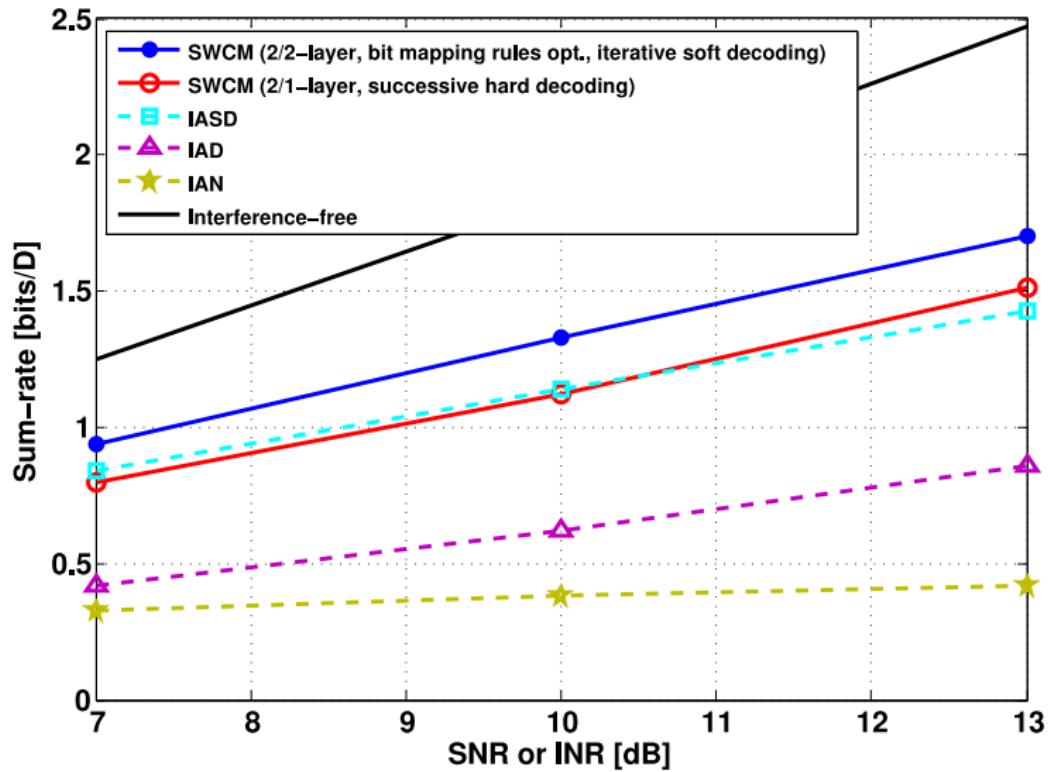
LTE turbo code / ≤ 8 -iteration LOG-MAP decoding at $b = 20$, $n = 660$ (13200 REs), BLER = 0.1

OFDM/Ped-B channel performance (Kim et al. 2015)



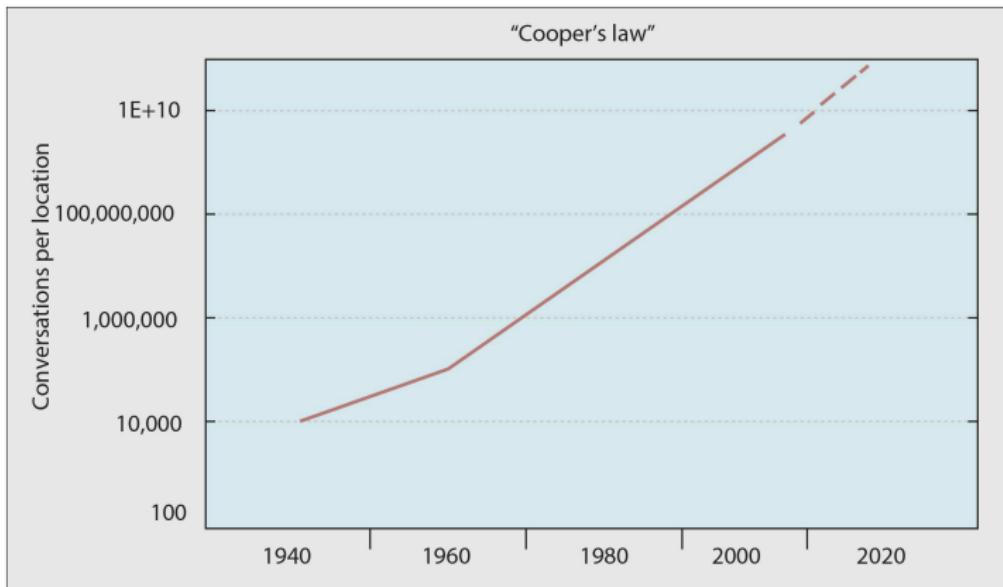
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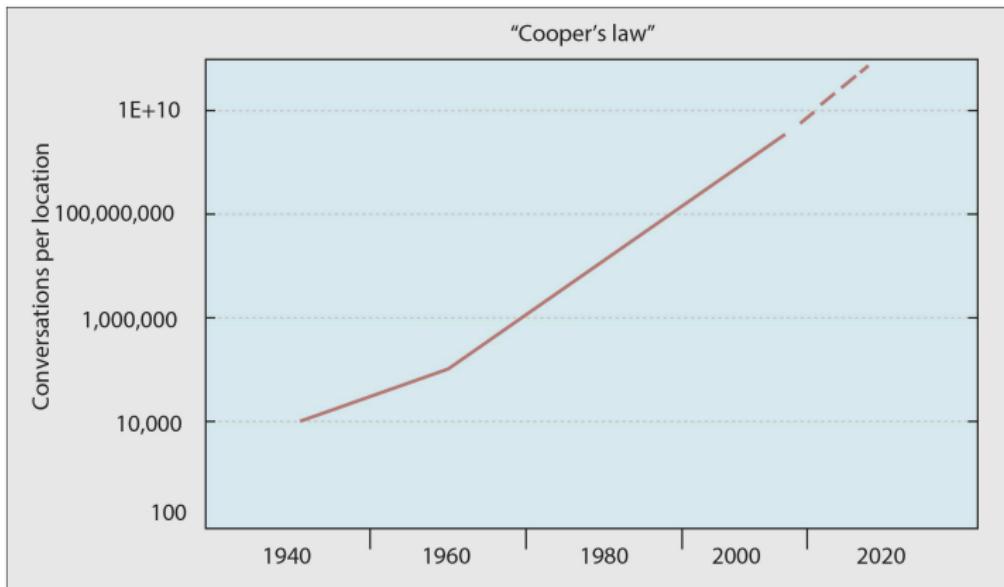
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Cooper's Law



Source: Arraycomm, Zander–Mähönen (2013)

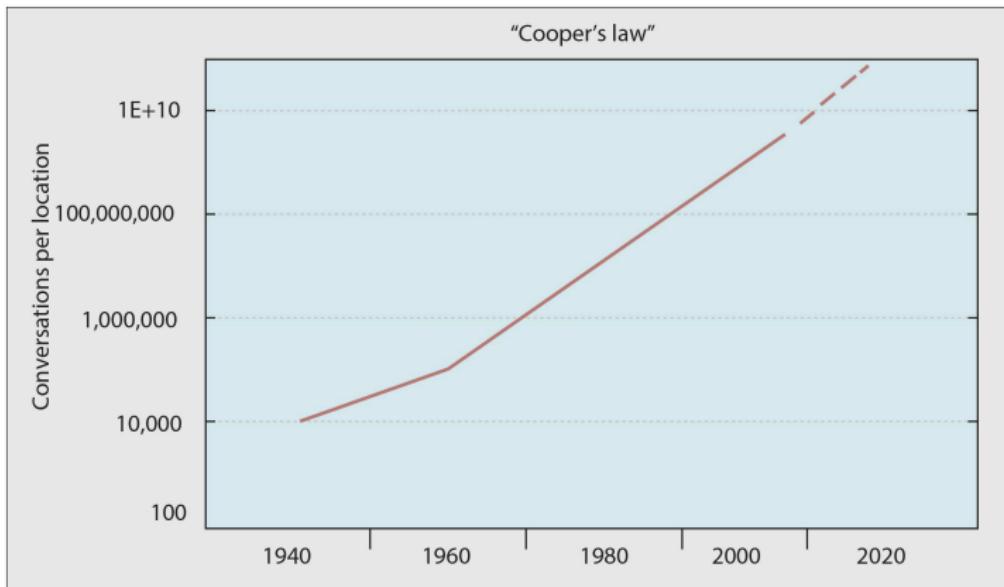
Cooper's Law



Source: Arraycomm, Zander–Mähönen (2013)

- Gain over the past 45 years = $10^6 \propto \eta W_{\text{sys}} N_{\text{BS}}$

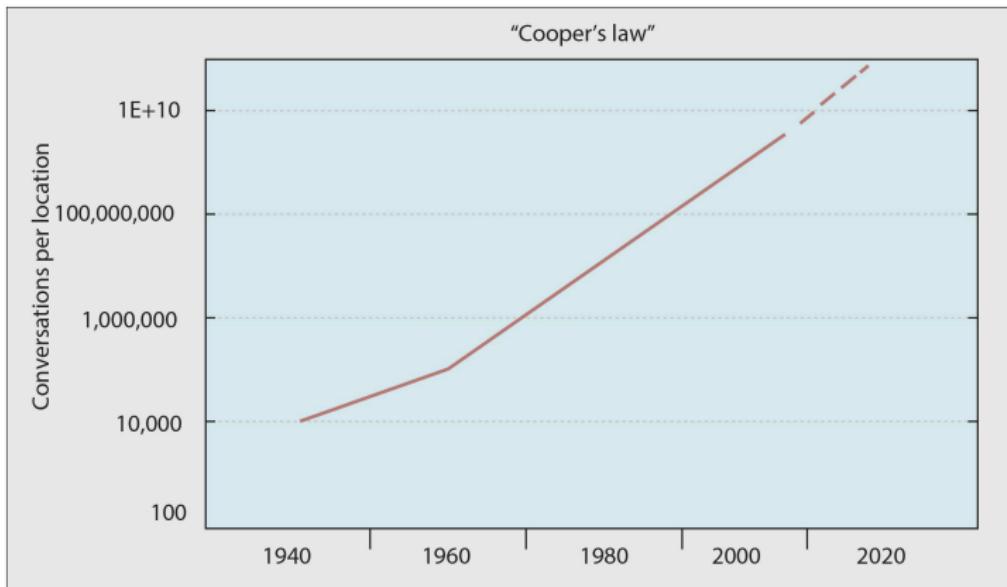
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Source: Arraycomm, Zander–Mähönen (2013)

- Gain over the past 45 years = $10^6 \propto \eta W_{\text{sys}} N_{\text{BS}}$
 - ▶ Spectral efficiency η : X 25

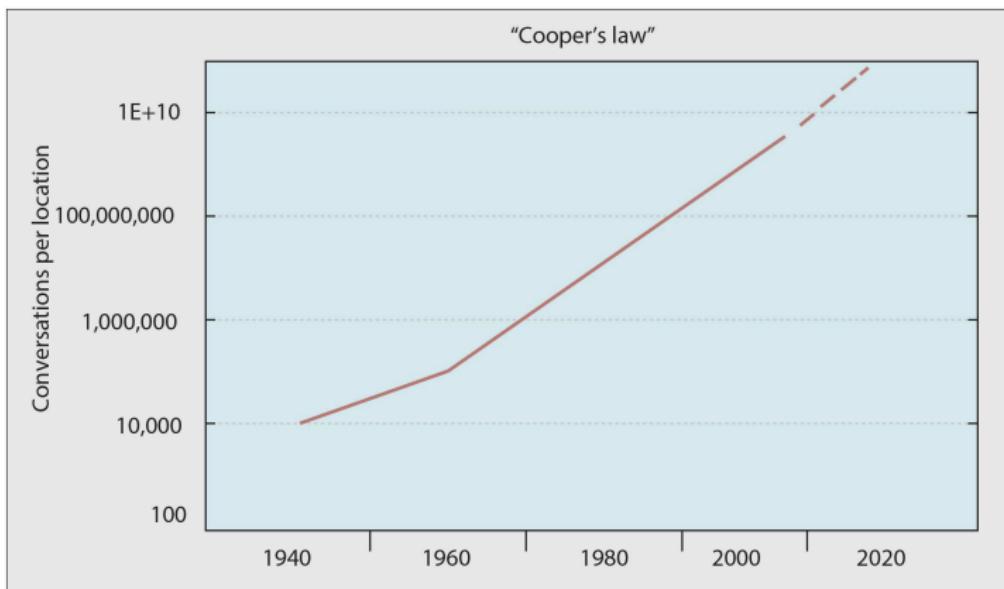
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- Gain over the past 45 years = $10^6 \propto \eta W_{\text{sys}} N_{\text{BS}}$
 - ▶ Spectral efficiency η : X 25
 - ▶ System bandwidth W_{sys} : X 25

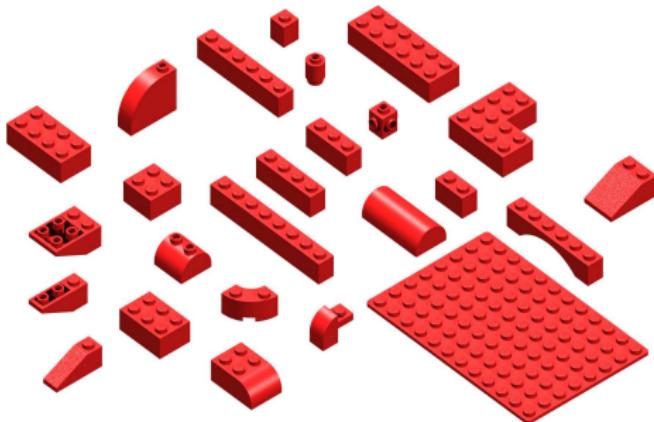
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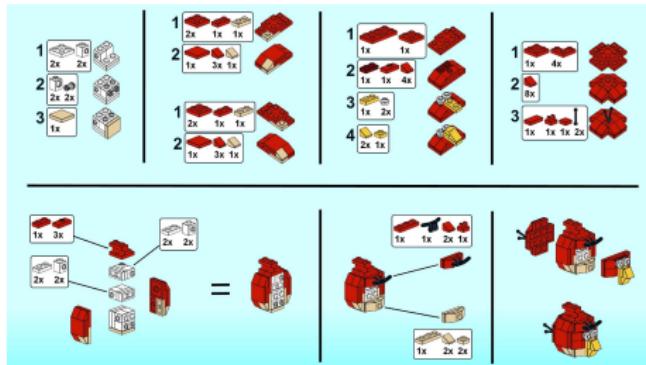
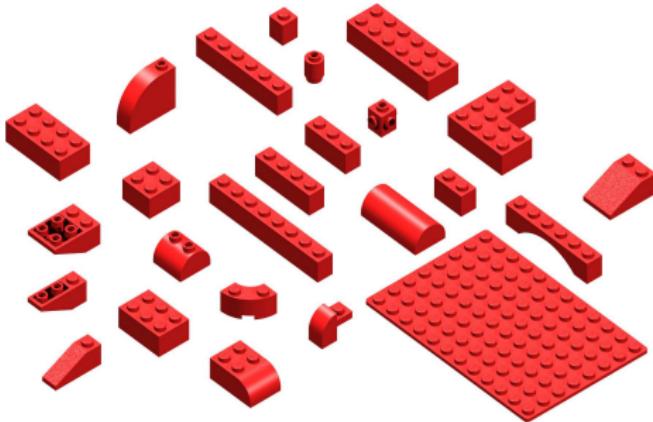
- Gain over the past 45 years = $10^6 \propto \eta W_{\text{sys}} N_{\text{BS}}$
 - ▶ Spectral efficiency η : X 25
 - ▶ System bandwidth W_{sys} : X 25
 - ▶ # of base stations N_{BS} : x 1600 (spatial reuse of frequency)

What's next?



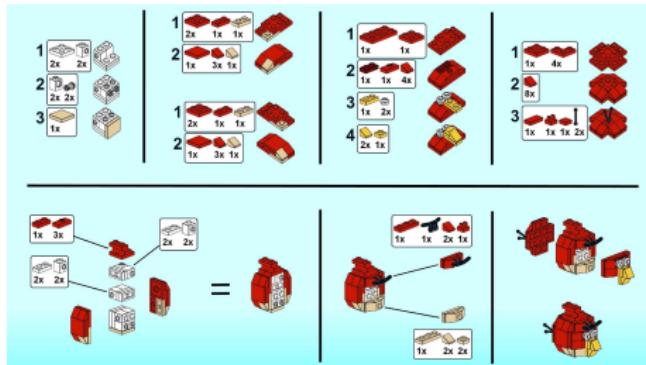
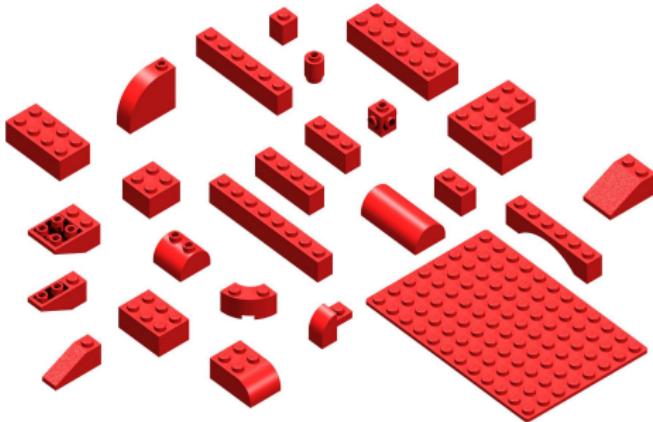
- Point-to-point codes (random coding)
- Superposition coding
- Successive cancellation decoding
- Simultaneous decoding
- Multicoding (writing on dirty paper)
- Random binning (Slepian–Wolf)

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- Noisy network coding
- Distributed decode-forward
- Index coding
- Distributed storage coding

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What's next?



Let's have fun building better networks!

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